



Network Coding

IN 3300

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- ❑ Network Coding Taxonomy
 - ❑ Reliable Multicast Transport
 - ❑ Forward Error Correction at Layer 2
 - ❑ Forward Error Correction at Application Layer
 - ❑ Examples from Research



Classification of Network Coding

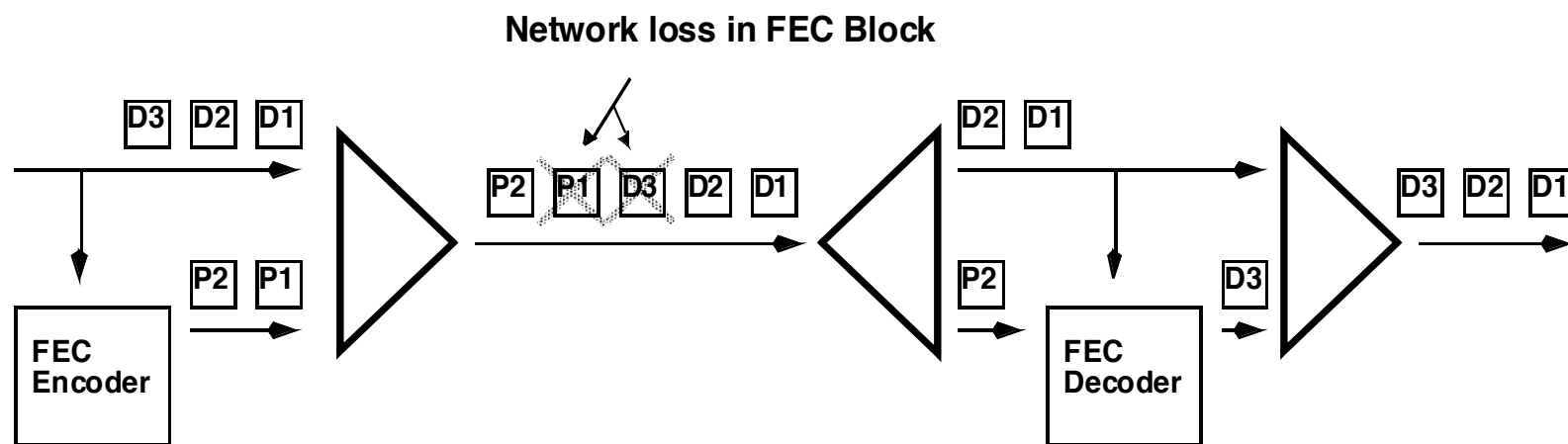
Starting with Forward Error Correction

Generalization to Network Coding



Packet-based Forward Error Correction (FEC)

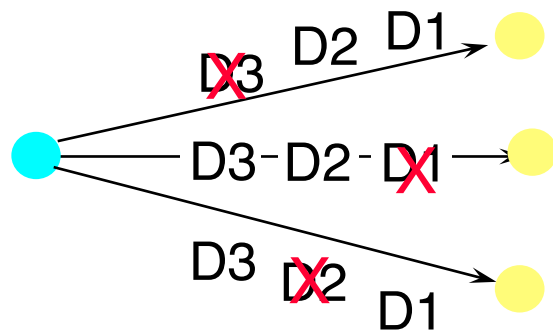
- ❑ k original data packets form a **Transmission Group (TG)**
- ❑ h parity packets derived from the k data packets
- ❑ any k received out of k+h are sufficient
- ❑ Assessment
 - + allows receiver (terminal) to recover lost packets
 - overhead at end-hosts
 - complex assessment of impact onto service quality
 - increased network load may increase loss probability



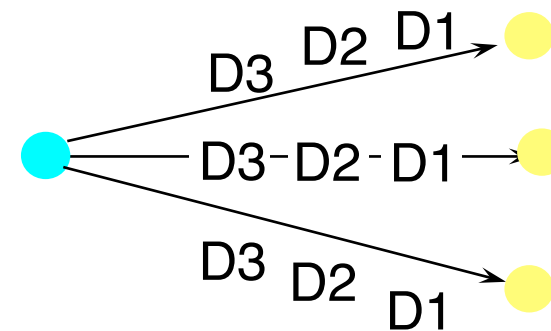


Potential Benefits of FEC

Initial Transmission

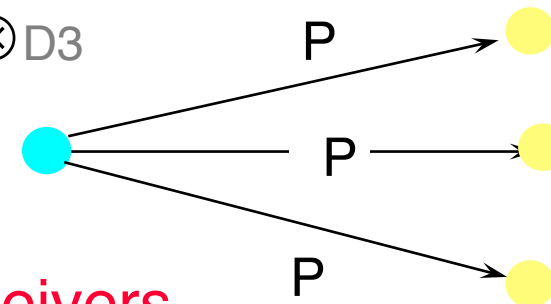


Data Retransmission



Parity Retransmission

$$P = D1 \otimes D2 \otimes D3$$



One parity packet can recover different data packets at different receivers



Network Coding

- Assumption behind traditional network traffic
 - Information is separate, although it may share network resources. (say, cars in highways or fluids in pipes).
 - Network coding breaks this assumption.

- Network Coding
 - A technology to combine several data packets into one or several output packets



- Core ideas:
 - Coded packet mixing improves
 - bottleneck traffic
 - Broadcast
 - Fountain approach simplifies scheduling/coordination
 - In wireless networks, opportunistic listening allows benefits for packets distributed and mixing over the air



- Intermediate nodes transmit packets that are functions of the received packets.

- Potential Benefits
 - throughput improvements in bottleneck scenarios
 - robustness to link failures
 - energy savings
 - simplified operation
 - etc.



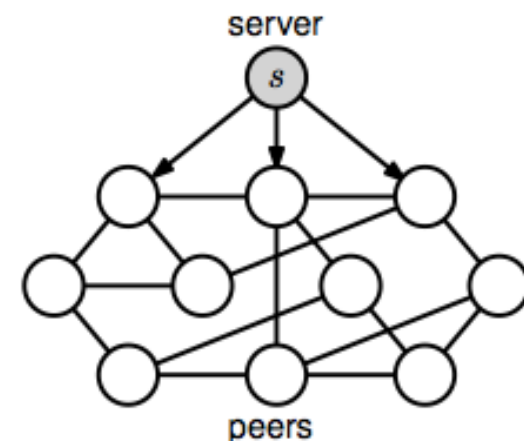
□ Applications

- Peer-to-peer multicast services
- Wireless Networks
 - network types:
sensor, adhoc, mobile, mesh, ...
 - functions:
routing in wireless networks
opportunistic listening
opportunistic forwarding



□ Multicast in Peer-to-Peer Networks

- Large application-level Service Data Unit (SDU) partitioned into blocks
- Service types
 - File distribution
 - On demand media streaming
 - Real-time media streaming
- Peers exchange blocks according to certain policy
 - Important policy: random blocks („random gossiping“)
- Property of solution
 - Scalable and robust to peers joining and leaving the system





- C. Gkantsidis and P. Rodriguez, “Network Coding for Large Scale Content Distribution”, IEEE INFOCOM 2005
 - “The performance benefits provided by network coding in terms of throughput can be more than 2-3 times better compared to transmitting unencoded blocks.”
 - <http://research.microsoft.com/pubs/67246/tr-2004-80.pdf>
- C. Gkantsidis, J. Miller, and P. Rodriguez, “Comprehensive view of a live network coding P2P system”, ACM IMC 2006
 - “Network coding incurs little overhead, both in terms of CPU processing and I/O activity, and it results in smooth, fast downloads, and efficient server utilization.”
 - <http://research.microsoft.com/pubs/69452/imc06.pdf>



Reliable Multicast Transport





Many Uses of Multicasting

- ❑ Teleconferencing
 - ❑ Distributed Games
 - ❑ Software/File Distribution
 - ❑ Video Distribution
 - ❑ Replicated Database Updates
- ⇒ multicast transport is done differently for each application



Multicast Application Modes

- ❑ Point-to-Multipoint:
Single Source, Multiple Receivers

- ❑ Multipoint-to-Multipoint:
Multiple Sources, Multiple Receivers

- ❑ Sources are receivers

- ❑ Sources are not receivers



Classification of Multicast Applications

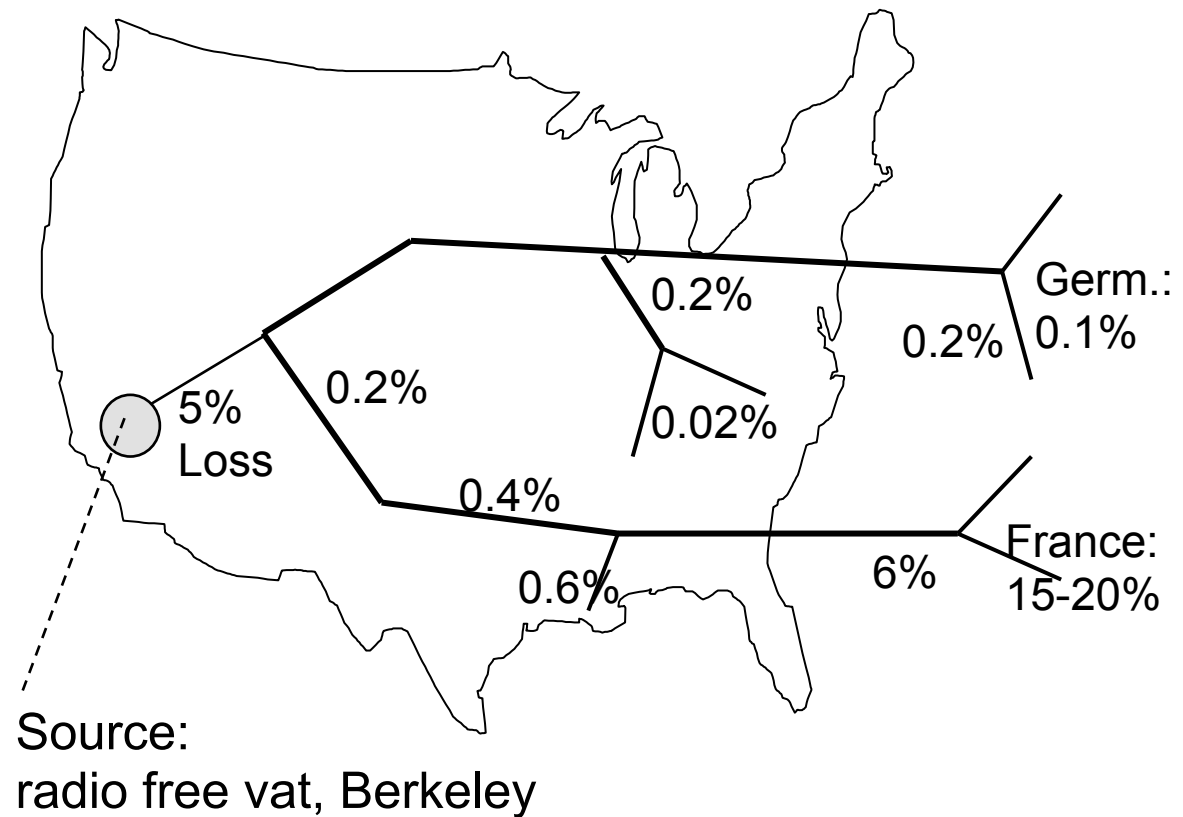
Transport service type	<i>Fully reliable multicast</i>	<i>Real-time multicast</i>
<i>Single source: 1:N</i>	Multicast-FTP; Software update	Audio-visual conference; Continuous Media Dissemination
<i>Multiple Sources M:N</i>	CSCW; Distributed computing	DIS; VR

- CSCW: Computer Supported Cooperative Work
- DIS: Distributed Interactive Simulation
- VR: Virtual Reality



Where Does Multicast Loss Occur

- Example measurements
(April 96, Yajnik, Kurose, Towsely, Univ. Mass., Amherst)





Simultaneous Packet Loss

- ❑ Q: distribution of number of receivers losing packet?

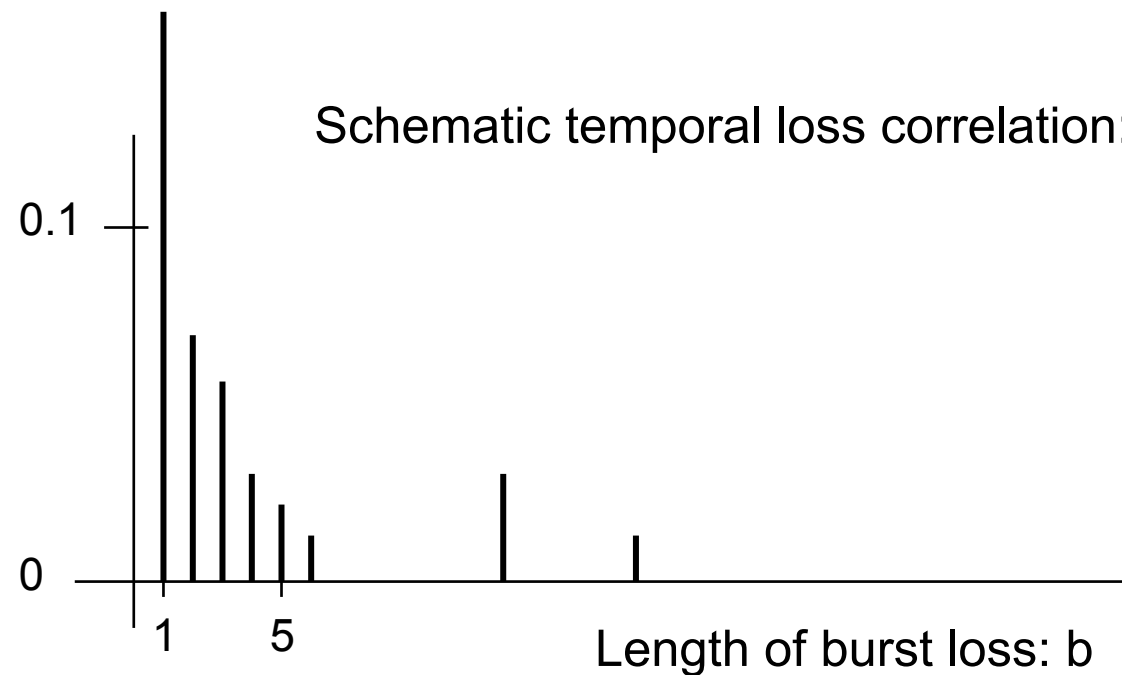
- ❑ Example dataset:
 - 47% packets lost somewhere
 - 5% shared loss
- ❑ Similar results across different datasets
- ❑ Models of packet loss (for protocol design, simulation, analysis):
 - star: end-end loss independently
 - full topology: measured per link loss independently
 - modified star: source-to-backbone plus star
 - ⇒ good fit for example data set



Q: do losses occur individually or in “bursts”?

- occasional long periods of 100% loss
- generally isolated losses
- occasional longer bursts

Prob. for burst
of length b





Reliable Multicast Challenge

- How to transfer data reliably from source to R receivers
- scalability: 10s - 100s - 1000s - 10000s - 100000s of receivers
- heterogeneity
 - different capabilities of receivers (processing power, buffer, protocol capabilities)
 - different network conditions for receivers (bottleneck bandwidths, loss rates, delay)
- feedback implosion problem



- Who retransmits
 - source
 - network / servers
 - other group member.
- Who detects loss
 - sender based: waiting for all ACKs
 - receiver based:
NAK, more receivers \Rightarrow faster loss detection.
- How to retransmit
 - Unicast
 - Multicast
 - Subgroup-multicast



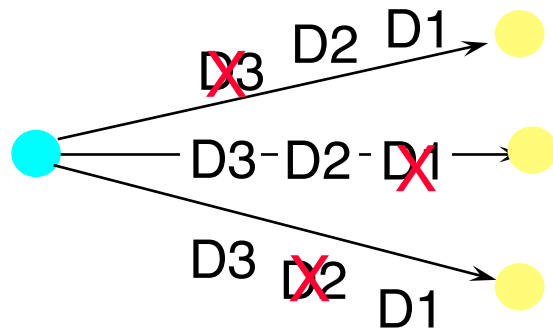
Approaches

- ❑ shift responsibilities to receivers (in contrast to TCP: sender is responsible for large share of functionality)
- ❑ feedback suppression (some feedback is usually required)
- ❑ multiple multicast groups (e.g. for heterogeneity problems; can be used statically or dynamically)
- ❑ local recovery (can be used to reduce resource cost and latency)
- ❑ server-based recovery
- ❑ forward error correction (FEC)
 - FEC for unicast: frequently no particular gain
 - FEC for multicast: gain may be tremendous!

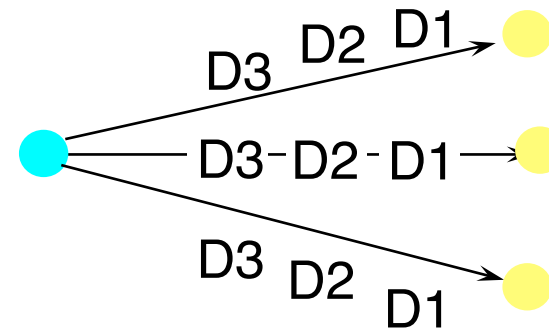


Multicast with FEC Error Correction

Initial Transmission

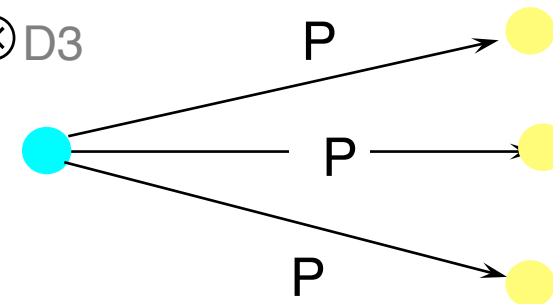


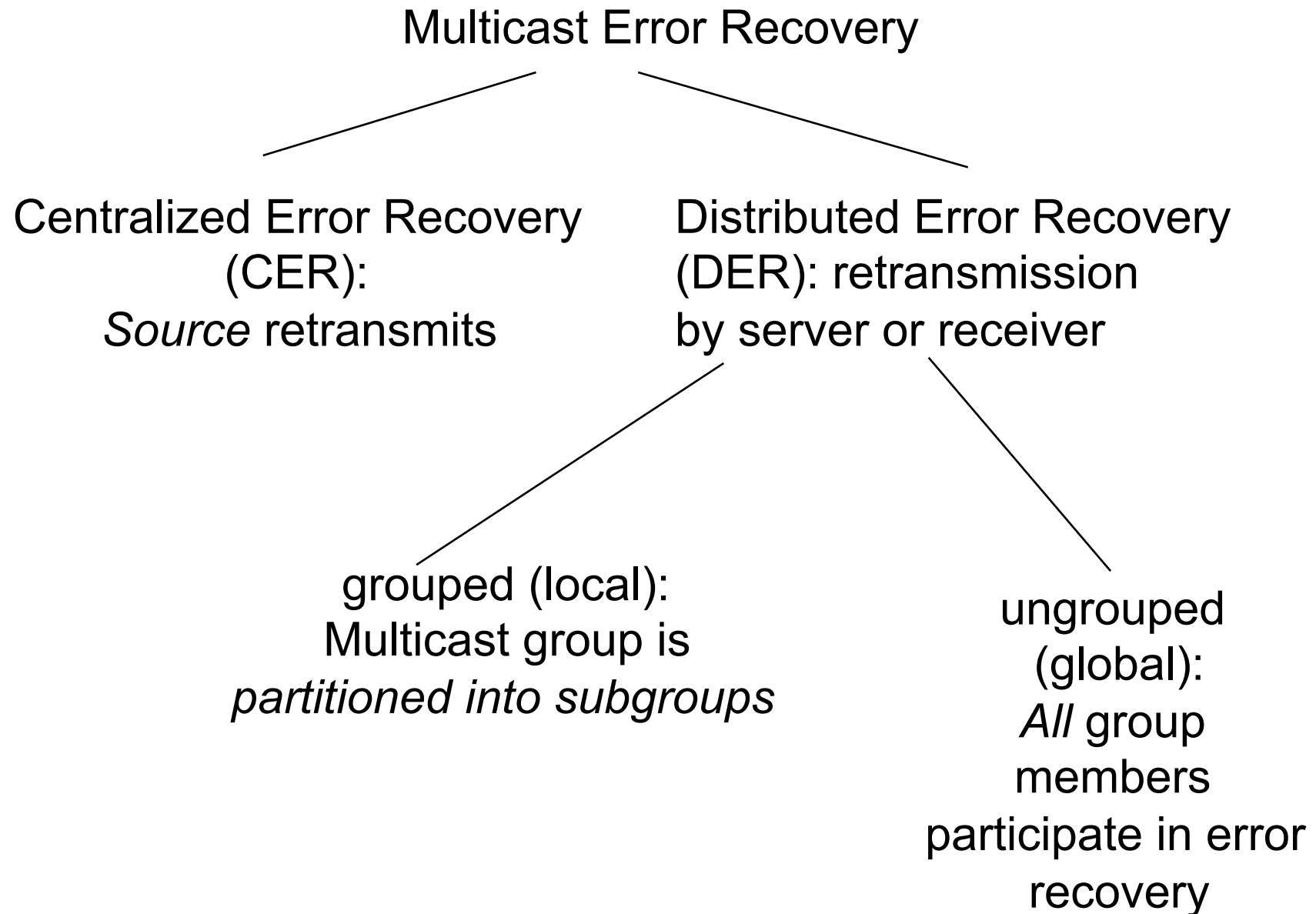
Data Retransmission



Parity Retransmission

$$P = D1 \otimes D2 \otimes D3$$







- Elements from Unicast:
 - Loss detection
 - Sender-based (ACK): 1 ACK per receiver and per packet; Sender needs a table of per-receiver ACK
 - Receiver-based (NAK): distributed over receivers; potentially only 1 NAK per lost packet
 - Loss recovery: ARQ vs. FEC
- Additional new elements for Multicast:
 - Mechanisms for control message **Implosion Avoidance**
 - Mechanisms to deal with *heterogeneous receivers*