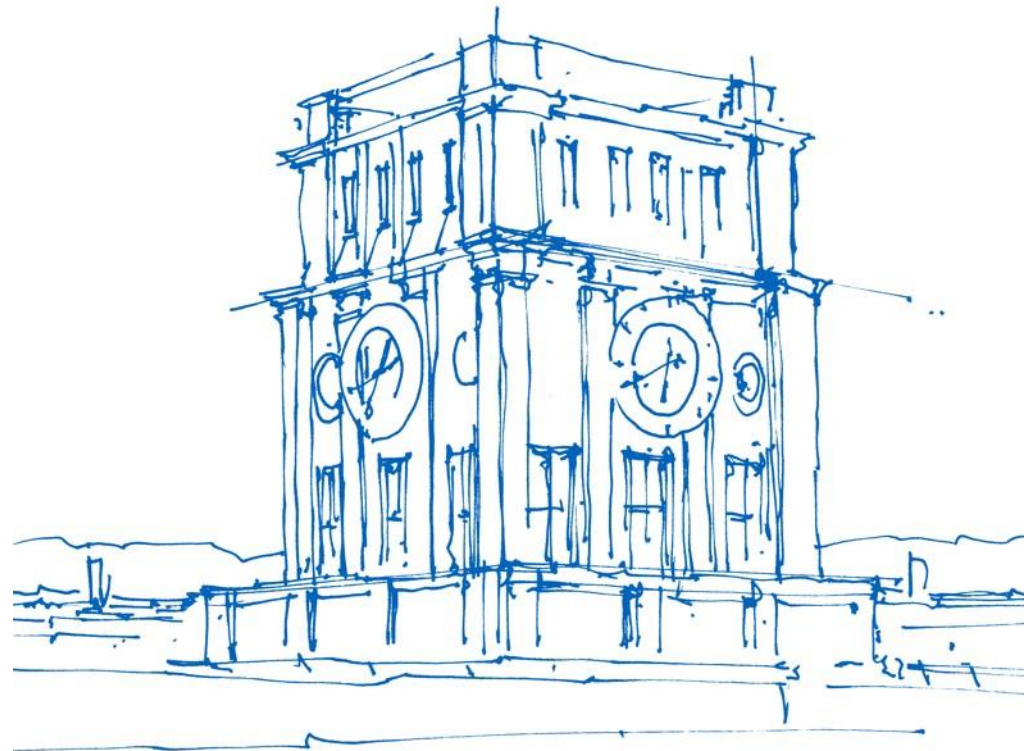


Firewall offloading based on SDN and NFV

ITG 5.2.2/5.2.4 05.12.2016

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Uhrenturm der TUM

Overview

- Introduction
 - Motivation
 - Main security requirements
- Network Function Virtualization
 - State of the Art
 - SDN & NFV Architecture
- Offloading Approach
 - ByteLim Logic
- Conclusion and open Questions

Software Defined Networking:

Central control improves programmability and makes innovations easier

Network Function Virtualization:

Run Network Functions on Commodity Hardware and in the cloud

- ✓ Reduce costs
- ✓ Use available resources flexible

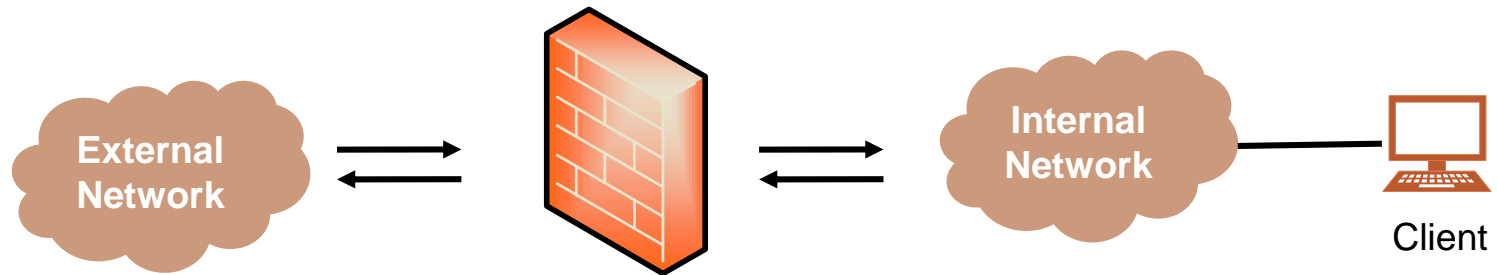
- How can we guarantee a certain security level in these environments?
- How can security related network functions be virtualized?

Main Security Requirements

- Isolation of services
 - Authentication and Authorization of devices/users/services
 - Isolation of flows
 - Stateless firewalling
- Stateful Firewalling
 - Check states of protocols
 - e.g. TCP, SIP
 - Normalization
 - e.g. filter non-standard DNS replies, filter html

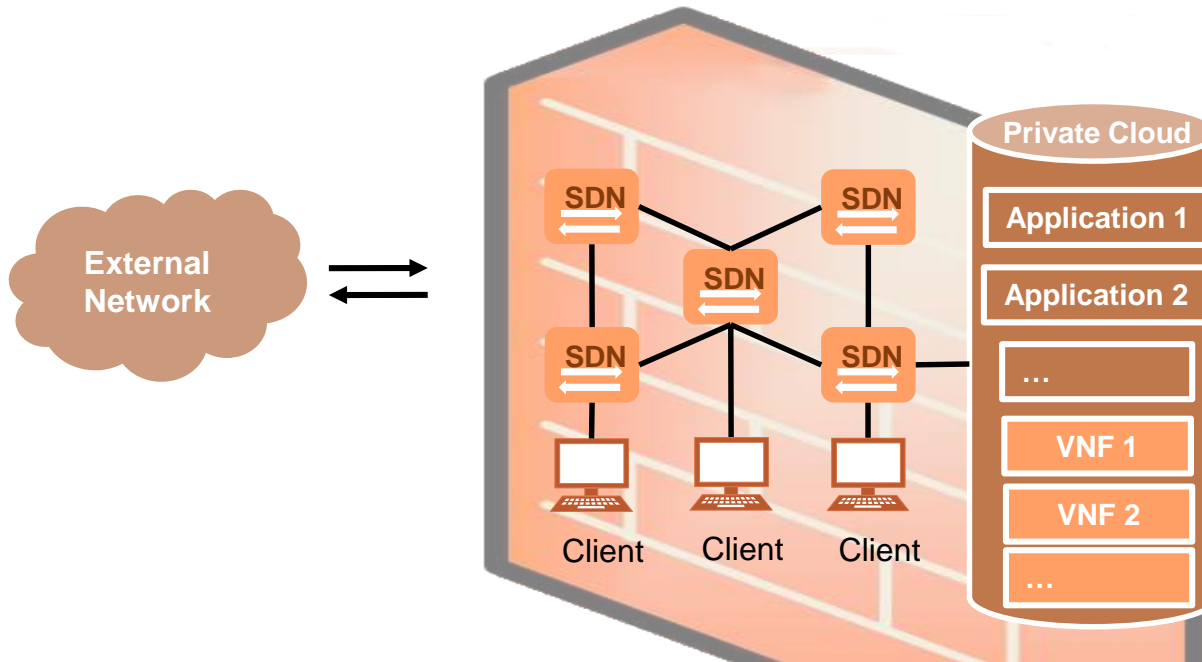
Stateless < Stateful < Application Layer

State of the Art Firewalls



- Firewall Resides on the Networks' Edge
- Control Plane (State) and Forwarding Plane not decoupled

SDN and NFV Network Security Architecture



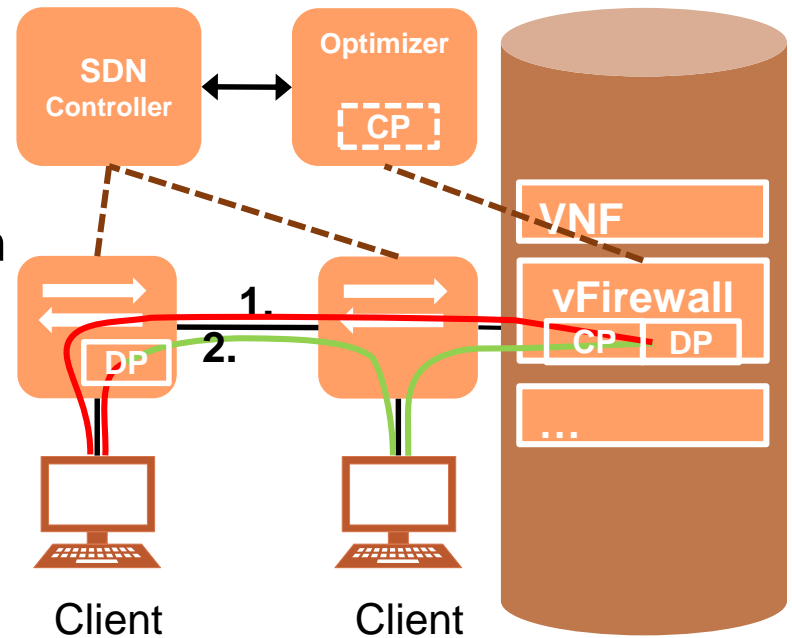
No distinct edge of the network

- Firewall has to filter „everywhere“
- Higher Load on Firewalls
- Potentially higher security

Offloading Approach

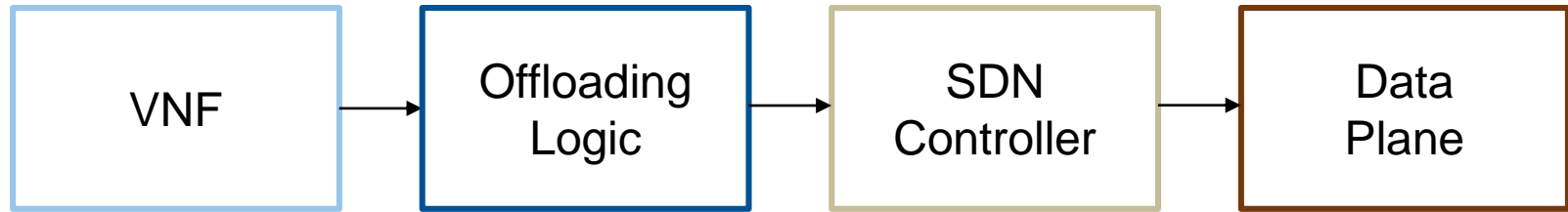
Combine NFV and SDN

- Traffic steering with SDN
- Some parts of the Network function with SDN
- Some parts offloaded to NE
- More Complex
- + leverages benefits of both approaches



Example TCP:

1. Connection Setup using VNF
2. Established connection using Hardware



- VNF signals connection state to Logic
- Logic decides for Offloading Flow
- Command to SDN Controller
- SDN Controller installs necessary rules in Hardware

Challenges

- Flow setup in switch

Can cause duplicates, packet loss and (short time) connection interruption

- Lack of Hardware

Current OpenFlow switches handle header rewrite in software – very low throughput

- Offloading Decision:

Flow classification algorithms needed

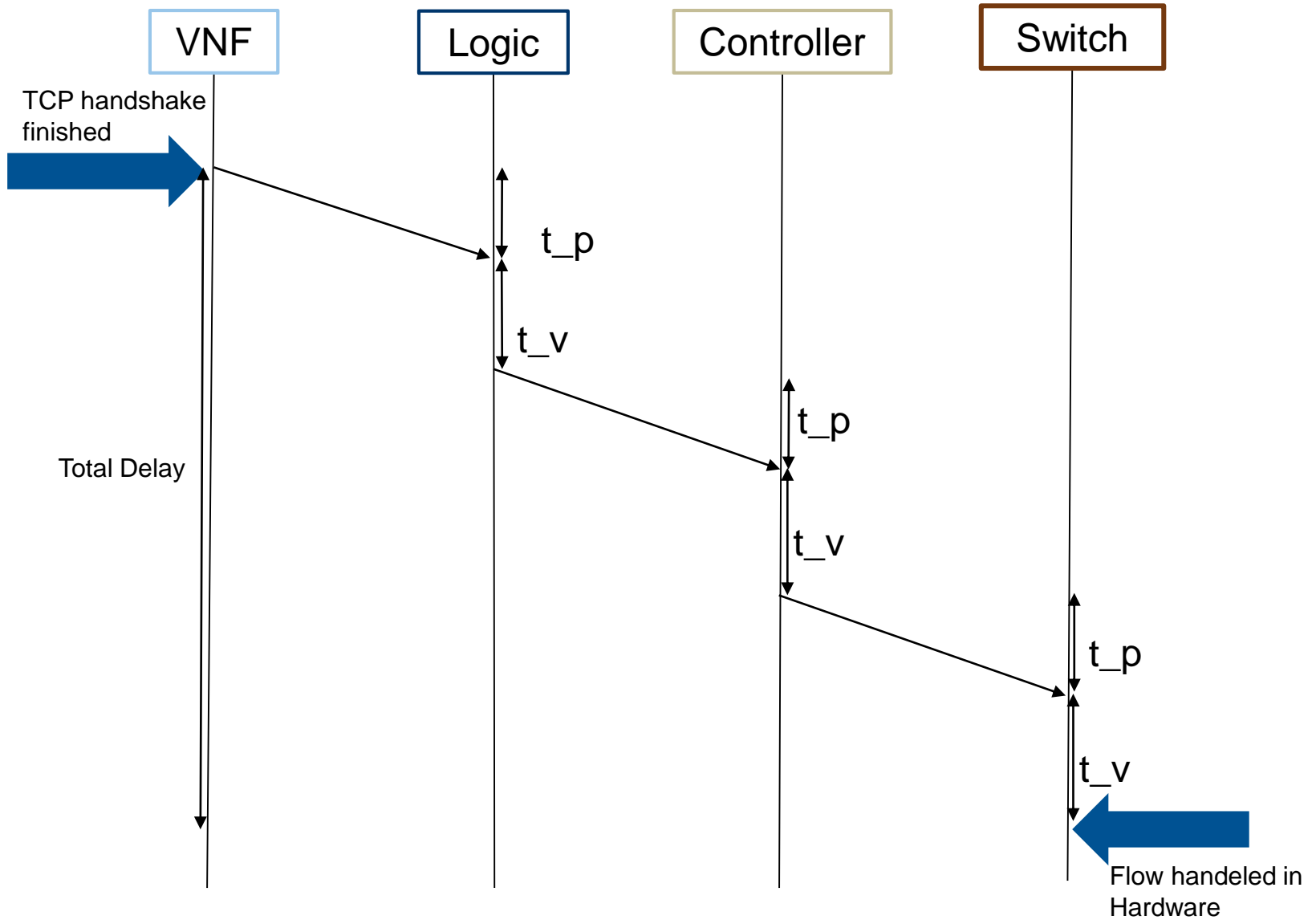
Which Flows can be offloaded and what are the gains?

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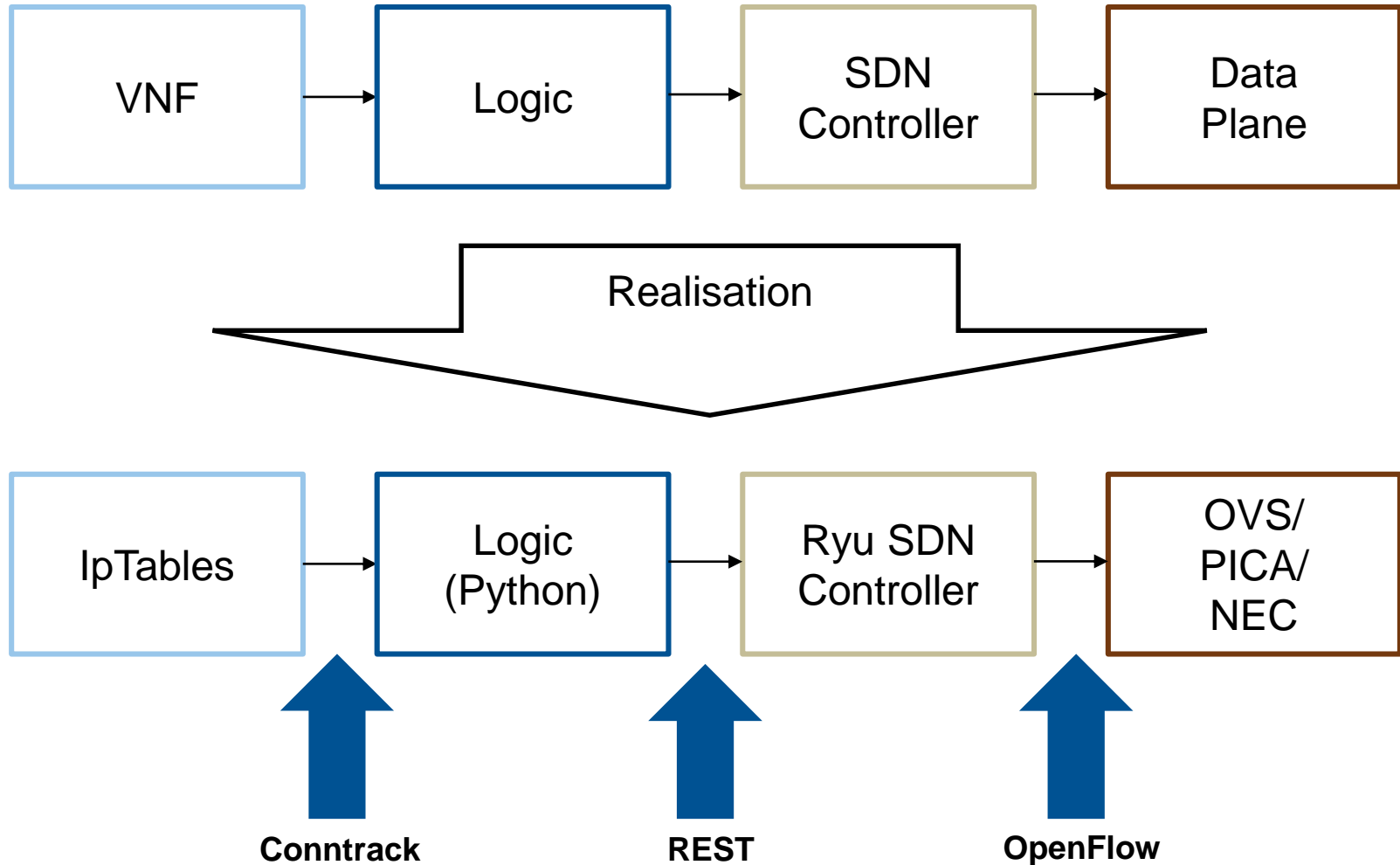
→ What limits the usage of Offloading?

- Flow capacity in the hardware
- Flow Setup Rate
- Delay from decision to active Offloading

Constraint: Delay in the complete system



Building Blocks Realisation



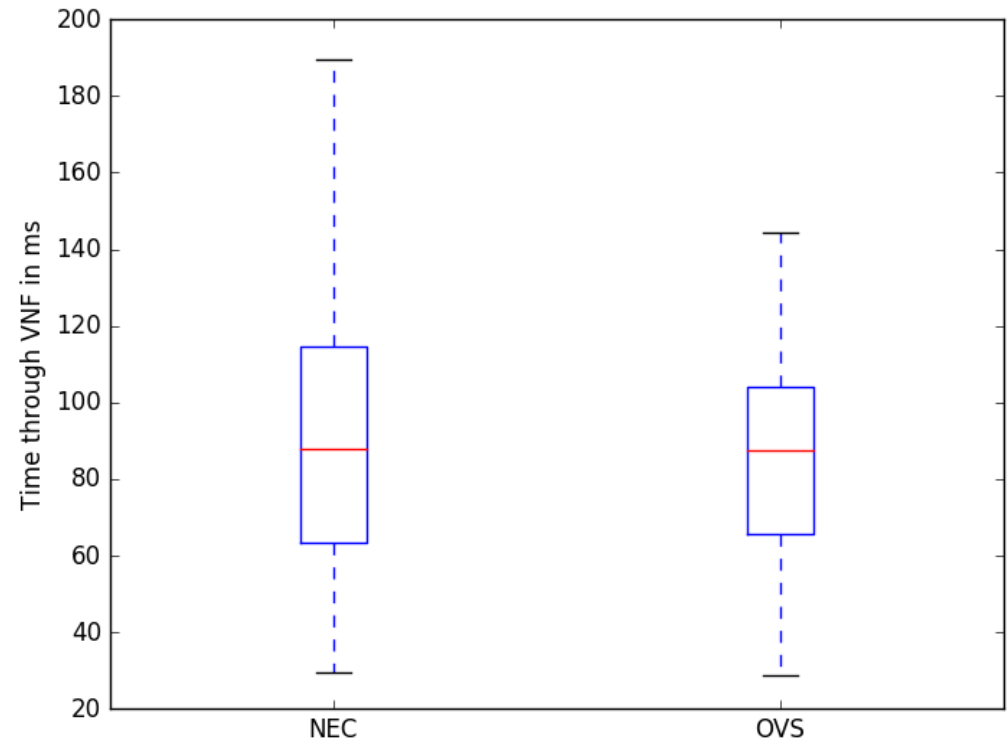
Total Delays

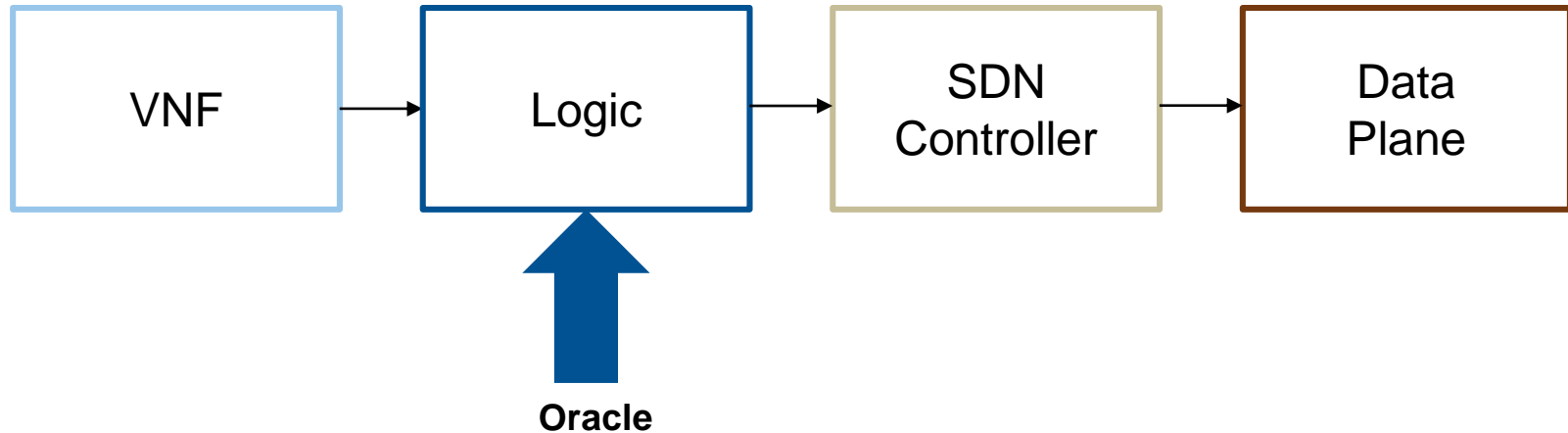
Results:

- ~ 100 ms for OVS & NEC

Implications:

- Not feasible for short lived flows like DNS
 - Delay > RTT
- Effects on TCP algorithm





- Which Flow metrics must be estimated by the Oracle?
- Gain i.e. cost of flow through VNF proportional to packet count

➔ Offload Flows with many Packets

Bytelimit Logic

- Logic decides based on Flowsize
 - Flows above a threshold are offloaded
- Oracle predicts Flowsize based on used application

Mathematical Description:

$f(x)$: PDF of flow size

x : Flow size

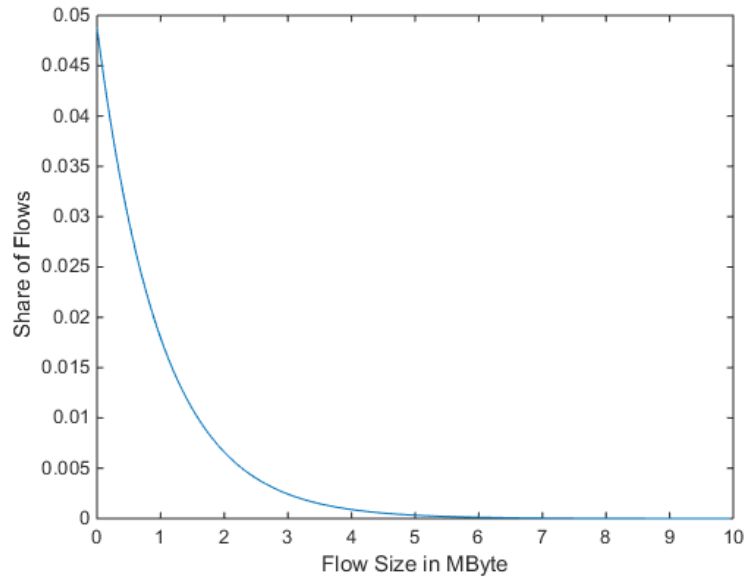
$P(x)$: PDF of Packets

$$P(x) = f(x) * x = \int f(x) \cdot x dx$$

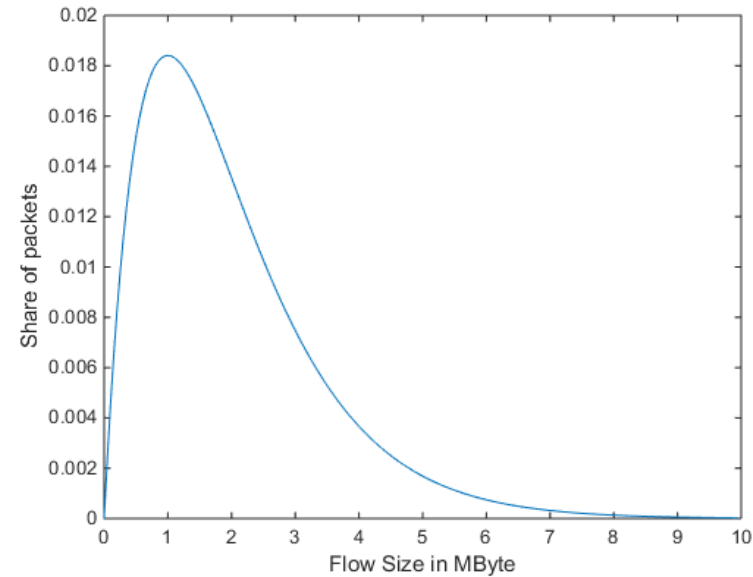
$$\xrightarrow{\text{discrete}} P = F \cdot X$$

Example: Negative Exponentially distributed Flowsize

Share of Flows $f(x)$



Share of Packets $P(x)$

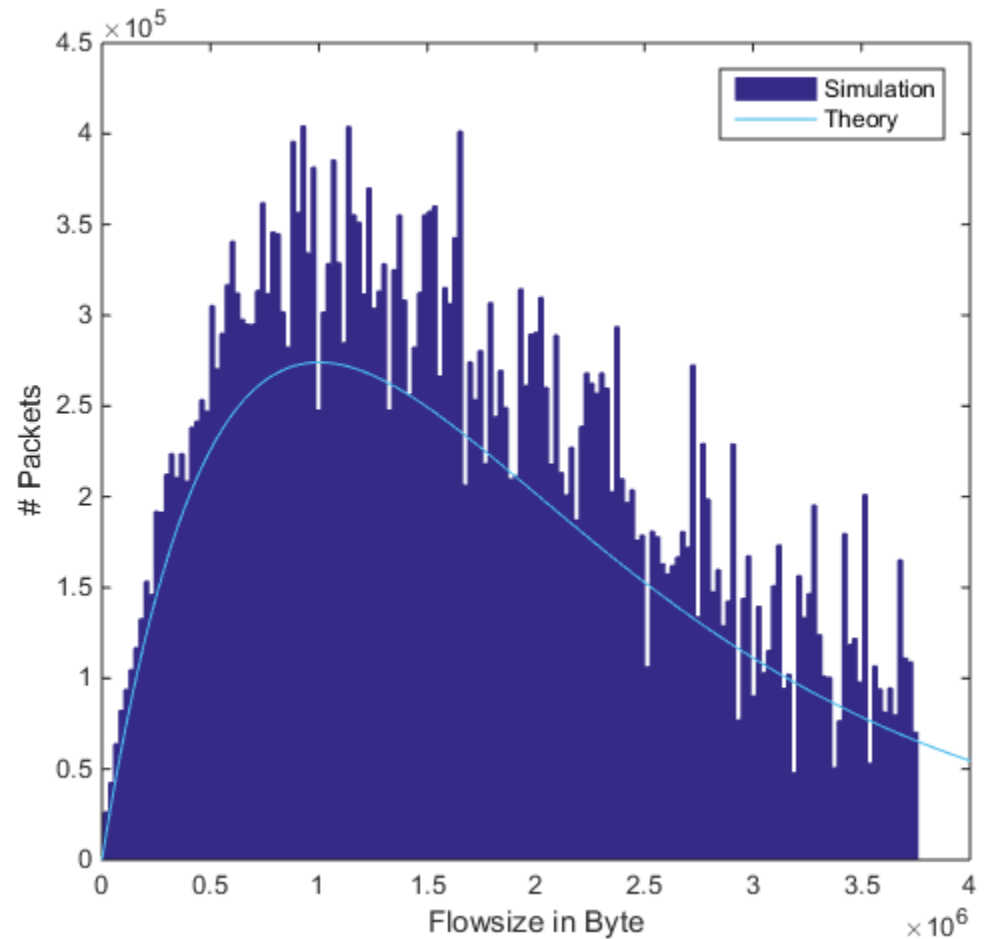


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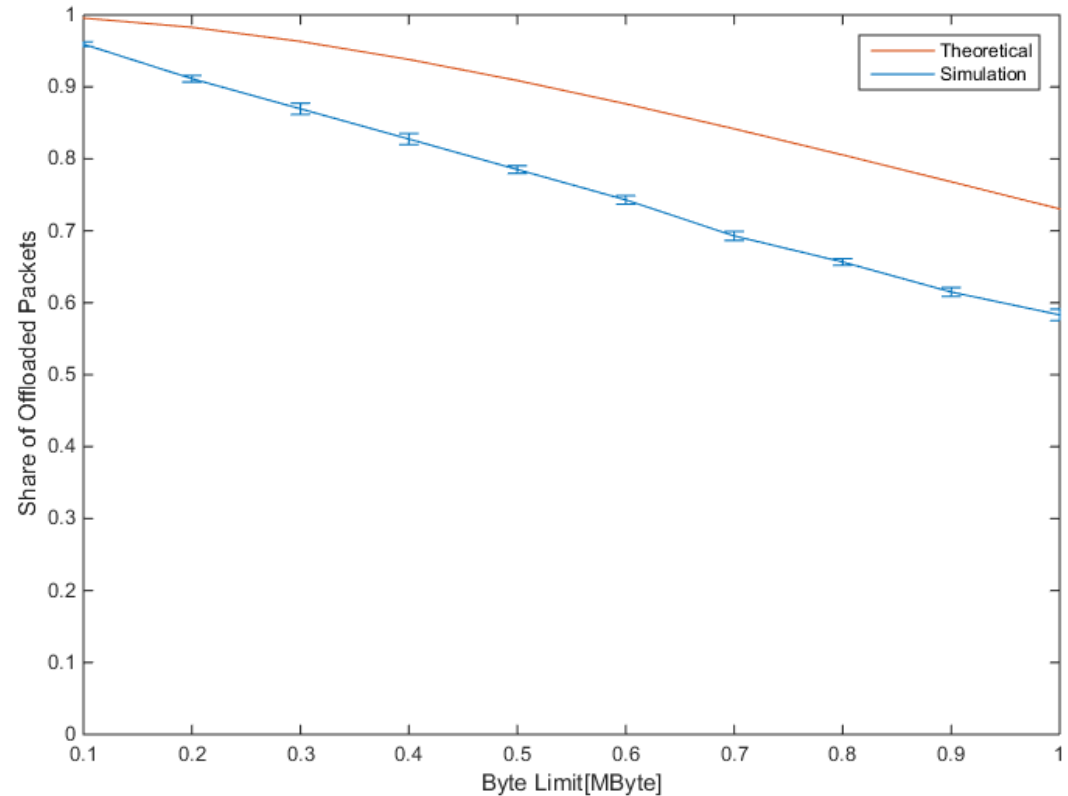
Simulation vs Theory

- Trendline similar
- Very big flows are hard to simulate



Simulation vs Theory

- Trendline similar
- Offset between simulation and theory
- Very big flows are hard to simulate



Conclusion and Open Questions

- Basic building blocks for offloading developed
- Bytlim logic shows promising results
 - Difference between simulation and theory should be evaluated

Open:

- How to predict Bytesize?
 - Simple classification by {source IP, Port}



Questions?