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# **Master Course Computer Networks IN2097**

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# Chapter: Quality of Service Support



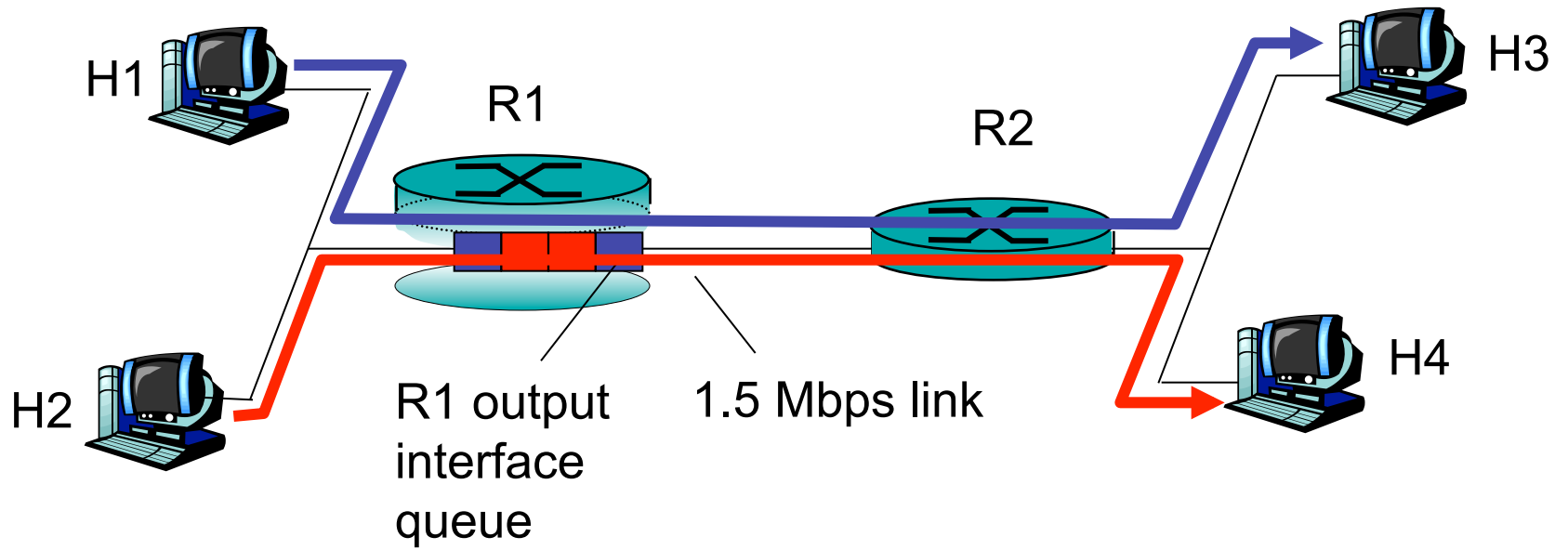


## Chapter outline – Quality-of-Service Support

- Providing multiple classes of service
- Providing QoS guarantees
- Signalling for QoS



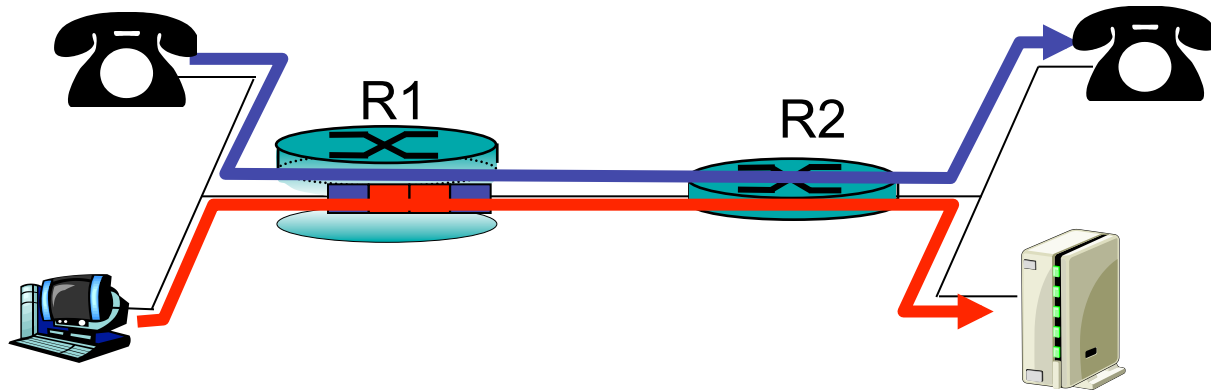
# Multiple Classes of Service: Scenario





## Scenario 1: mixed FTP and audio

- Example: 1Mbps IP phone, FTP or NFS share 1.5 Mbps link.
  - bursts of FTP or NFS can congest router, cause audio loss
  - want to give priority to audio over FTP



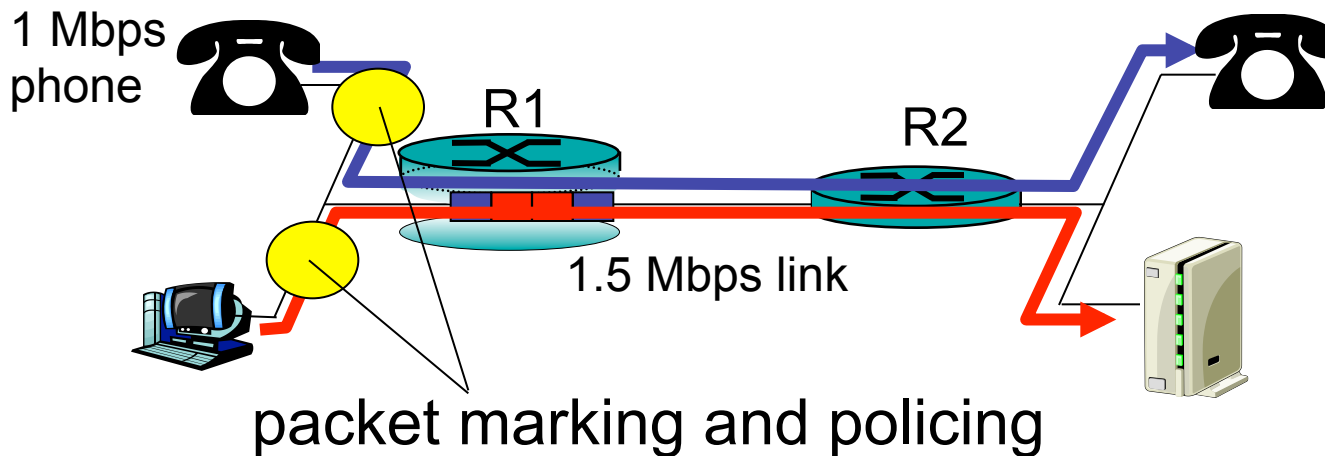
### Principle 1

packet marking needed for router to distinguish between different classes; and new router policy to treat packets accordingly



# Principles for QOS Guarantees (more)

- ❑ What if applications misbehave (audio sends higher than declared rate)
  - policing: force source adherence to bandwidth allocations
- ❑ Marking and policing at network edge:
  - similar to ATM UNI (User Network Interface)



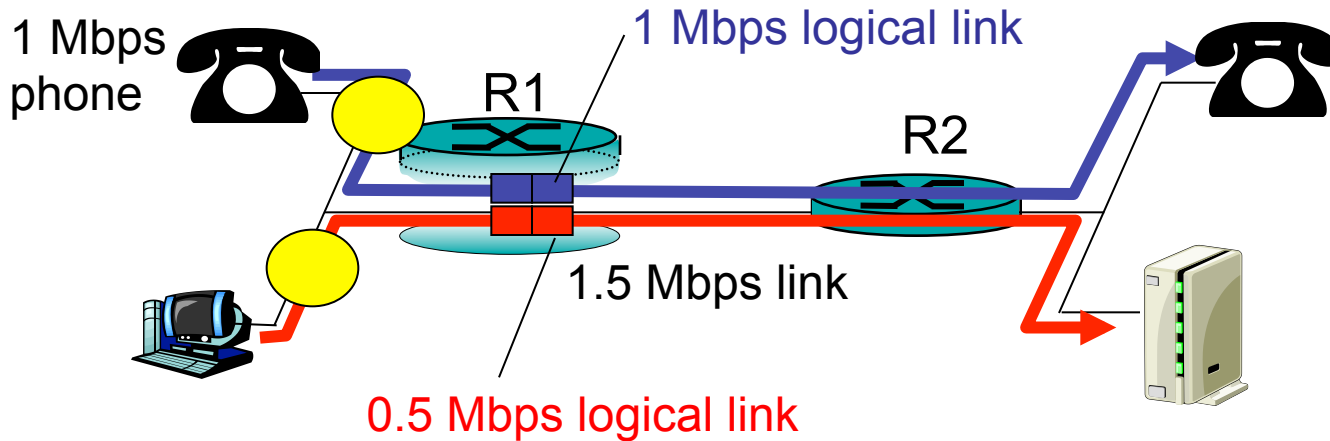
## Principle 2

provide protection (*isolation*) for one class from others



# Principles for QOS Guarantees (more)

- ❑ Allocating *fixed* (non-sharable) bandwidth to flow: *inefficient* use of bandwidth if flows doesn't use its allocation



## Principle 3

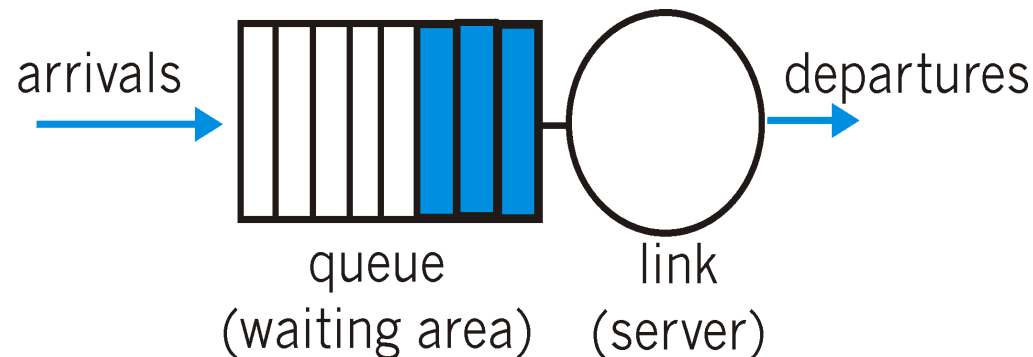
While providing **isolation**, it is desirable to use resources as efficiently as possible

- ⇒ allocate sharable bandwidth to logical link
- issue: sharing policy (scheduling, discarding) to be defined



# Scheduling And Policing Mechanisms

- ❑ **Scheduling:** choose next packet to send on link
- ❑ **FIFO (first in first out) scheduling:** send in order of arrival to queue
- ❑ **Discard policy:** if packet arrives to full queue: who to discard?
  - Tail drop: drop arriving packet
  - priority: drop/remove on priority basis
  - random: drop/remove randomly



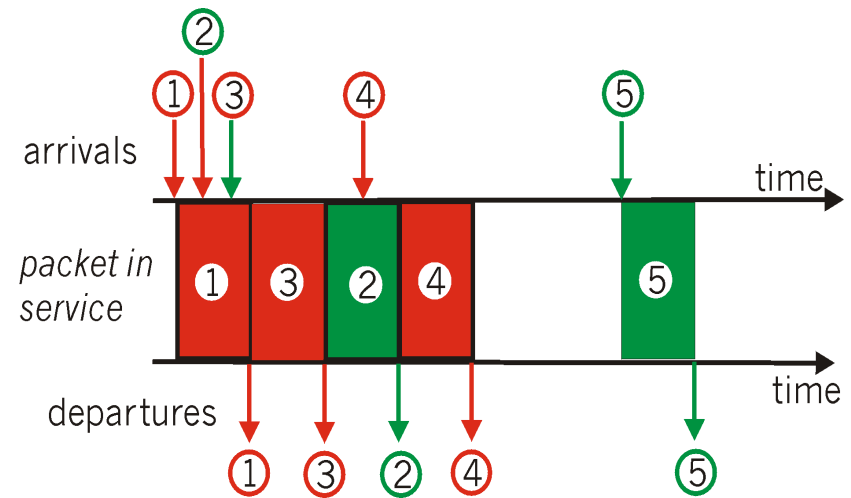
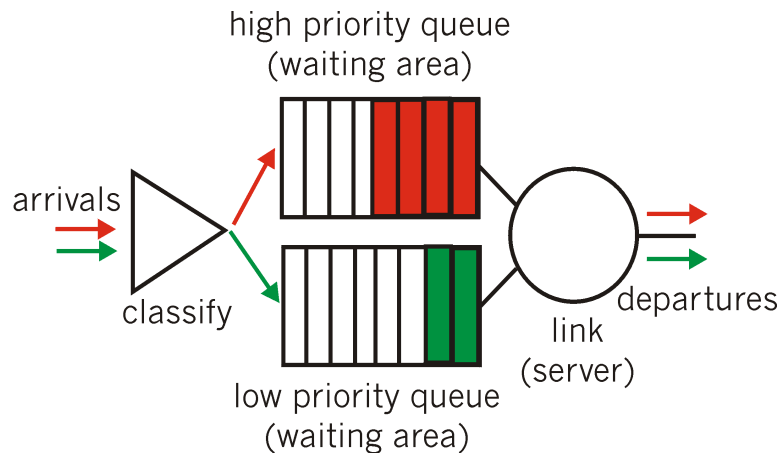




# Scheduling Policies: more

**Priority scheduling:** transmit highest priority queued packet

- multiple *classes*, with different priorities
  - class may depend on marking, or other header info, e.g. IP source/dest, port numbers, etc..

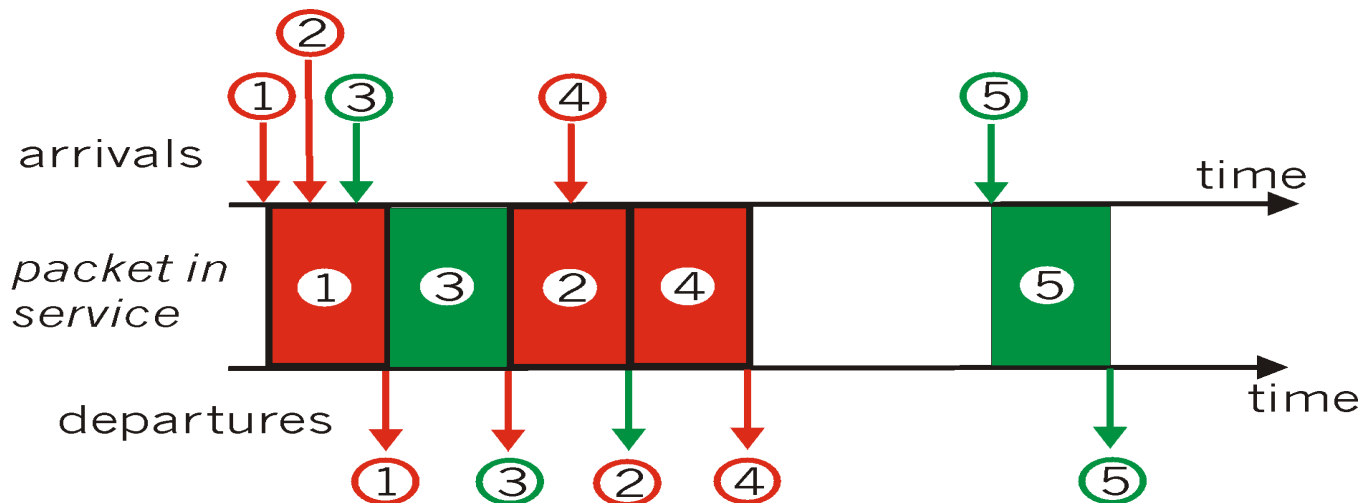




# Scheduling Policies: more

## Round robin scheduling:

- ❑ multiple classes
- ❑ cyclically scan class queues, serving one from each class (if available)

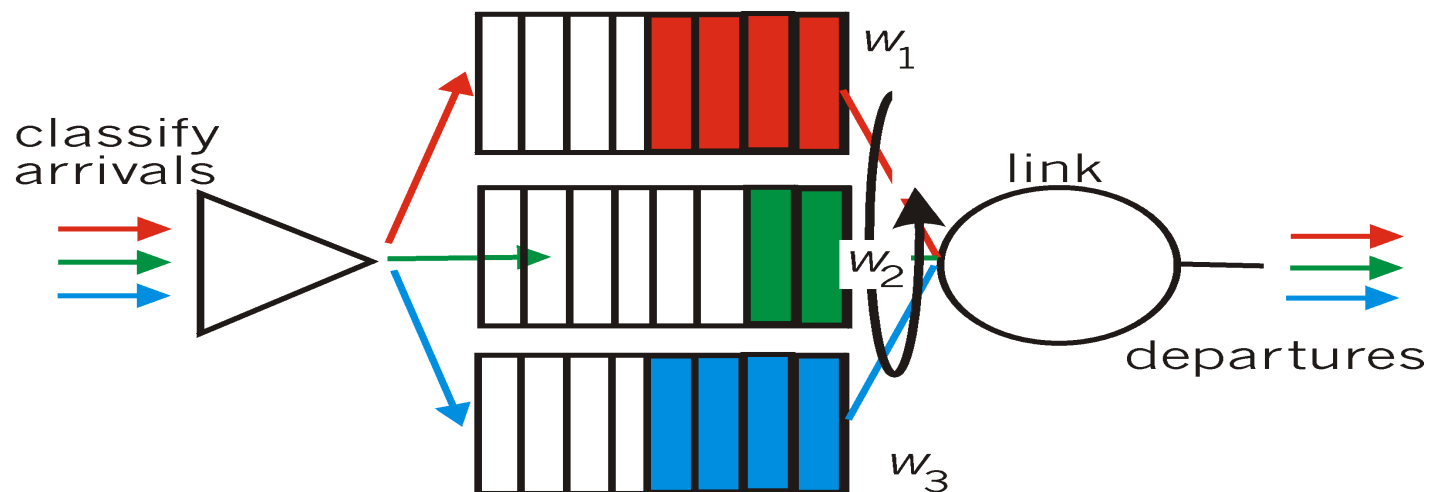




# Scheduling Policies: more

## Weighted Fair Queuing:

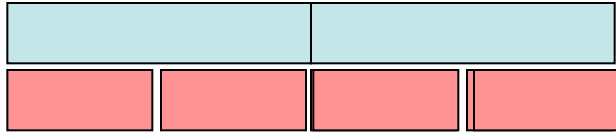
- ❑ each class gets weighted amount of service in each cycle
- ❑ when all classes have queued packets, class  $i$  will receive a bandwidth ratio of  $w_i / \sum w_j$  (for all  $j$  classes that have packets in queue)
- ❑ ill-behaved traffic classes only punish themselves
- ❑ Parekh and Galagher showed that combination with  $\rightarrow$  *leaky bucket policing* allows end-to-end delay bounds



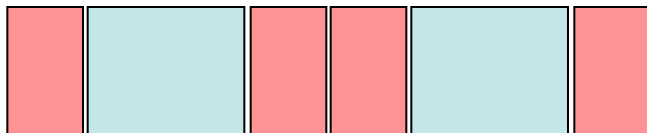


# WFQ and Packets

- ❑ Fluid Flow System (Processor Sharing)
  - work-conserving scheduling without scheduling overhead
  - fluid flow: conceptually bit-by-bit weighted round robin



- ❑ Packet-by-Packet scheduling
  - approach: use finishing time of packet in fluid system as priority for choosing next packet



- issue: arrival of packets of new flow  
⇒ virtual time (round number) finishing time



# Policing Mechanisms

Goal: limit traffic to not exceed declared parameters

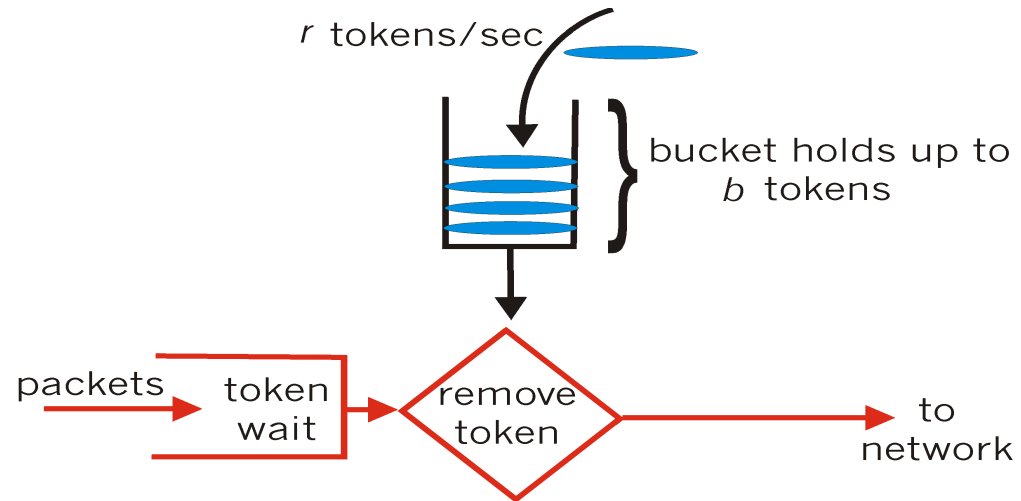
Three commonly used criteria:

- ❑ *(Long term) Average Rate:* how many packets can be sent per unit time (in the long run)
  - crucial question: what is the interval length:  
100 packets per sec  
or 6000 packets per min have same average!
- ❑ *Peak Rate:* e.g., 6000 packets per min (ppm) avg.;  
1500 pps (90000 ppm) peak rate
- ❑ *(Max.) Burst Size:* max. number of packets sent consecutively



# Policing Mechanisms

Token Bucket: limit input to specified Burst Size and Average Rate.

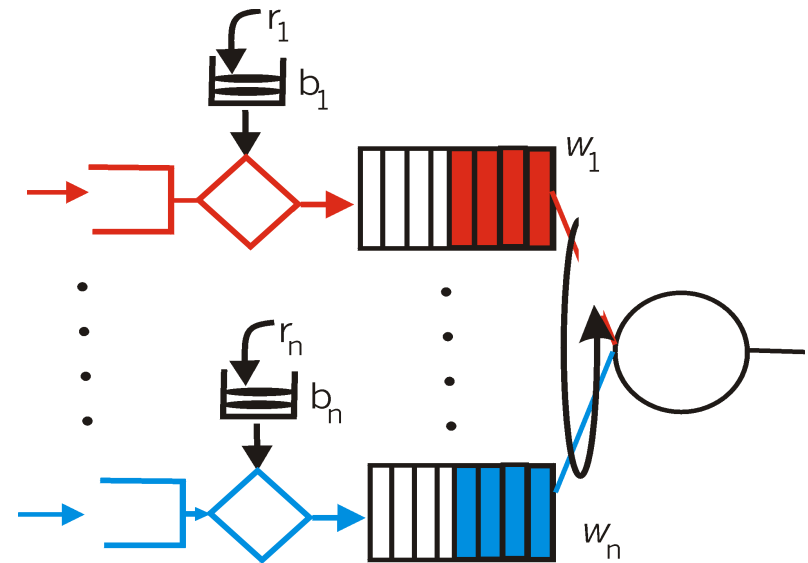
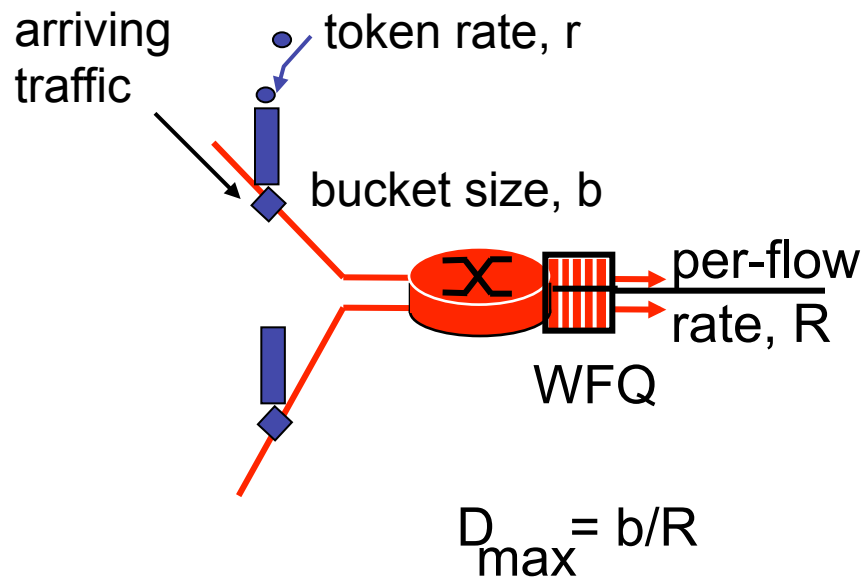


- ❑ bucket can hold  $b$  tokens  $\Rightarrow$  limits maximum burst size
- ❑ tokens generated at rate  $r$  token/sec unless bucket full
- ❑ *over interval of length  $t$ : number of packets admitted less than or equal to  $(r t + b)$ .*



# Policing Mechanisms (more)

- token bucket, WFQ combined provide guaranteed upper bound on delay, i.e., *QoS guarantee*





# IETF Differentiated Services

- Want “qualitative” service classes
  - “behaves like a wire”
  - relative service distinction: Platinum, Gold, Silver
- *Scalability*: simple functions in network core, relatively complex functions at edge routers (or hosts)
  - in contrast to IETF Integrated Services: signaling, maintaining per-flow router state difficult with large number of flows
- Don't define service classes, provide functional components to build service classes





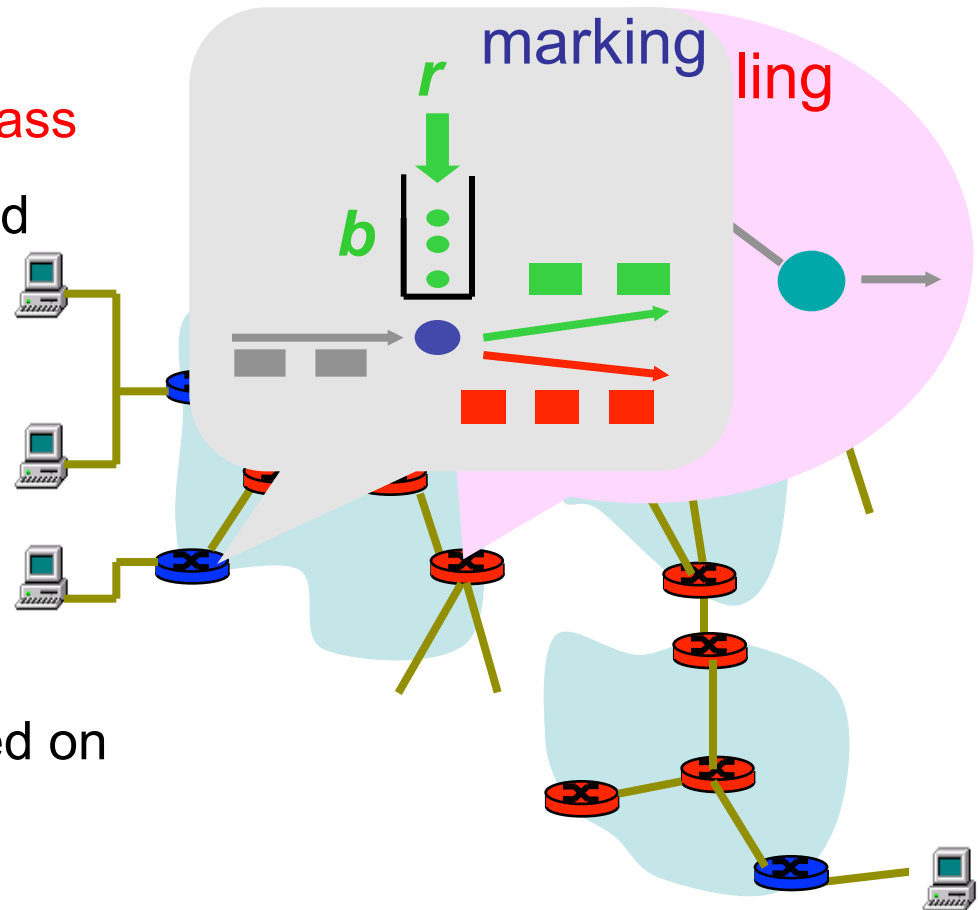
# Diffserv Architecture

## Edge router:

- ❑ per-flow traffic management
- ❑ marks packets according to **class**
- ❑ marks packets as **in-profile** and **out-profile**

## Core router:

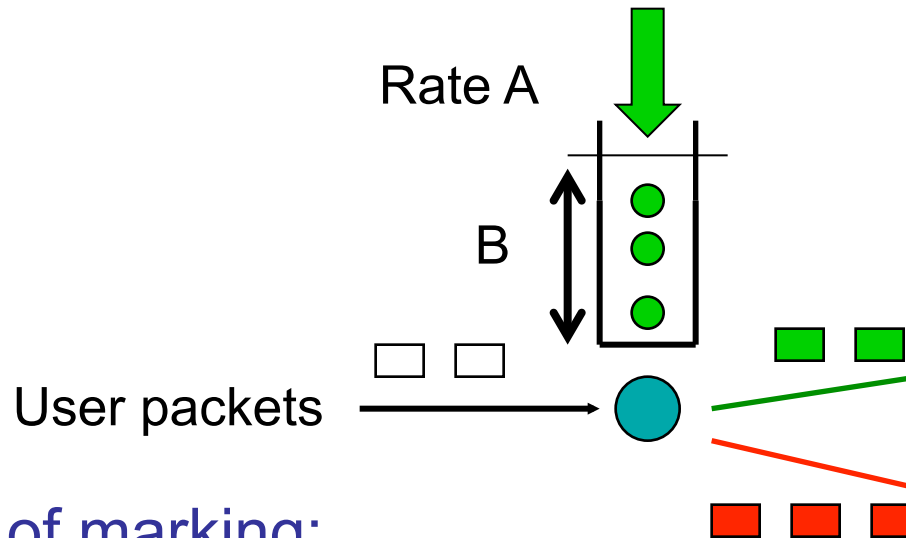
- ❑ **per class** traffic management
- ❑ buffering and scheduling based on **marking** at edge
- ❑ preference given to **in-profile** packets





# Edge-router Packet Marking

- **Profile:** pre-negotiated rate A, bucket size B
- Packet marking at edge based on **per-flow** profile



## Possible usage of marking:

- class-based marking: packets of different classes marked differently
- intra-class marking: conforming portion of flow marked differently than non-conforming one



## Classification and Conditioning

- ❑ Packet is marked in the Type of Service (TOS) in IPv4, and Traffic Class in IPv6
- ❑ 6 bits used for Differentiated Service Code Point (DSCP) and determine PHB that the packet will receive
- ❑ 2 bits can be used for congestion notification: Explicit Congestion Notification (ECN), RFC 3168

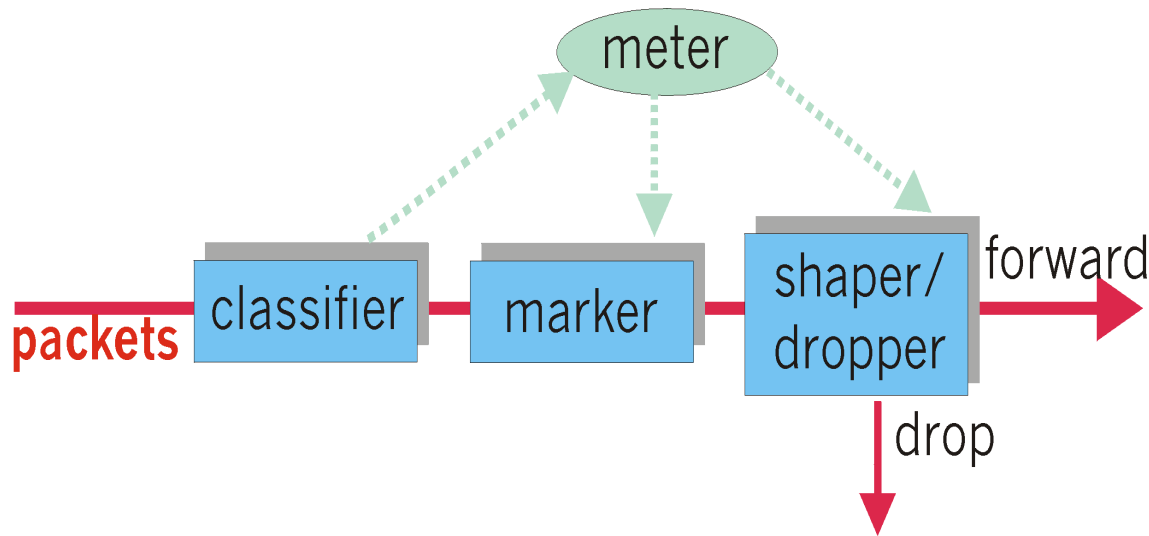




# Classification and Conditioning

May be desirable to limit traffic injection rate of some class:

- ❑ user declares traffic profile (e.g., rate, burst size)
- ❑ traffic metered, shaped or dropped if non-conforming





## Forwarding (PHB)

- ❑ PHB result in a different observable (measurable) forwarding performance behavior
- ❑ PHB does not specify what mechanisms to use to ensure required PHB performance behavior
- ❑ Examples:
  - Class A gets  $x\%$  of outgoing link bandwidth over time intervals of a specified length
  - Class A packets leave first before packets from class B



## Forwarding (PHB)

PHBs being developed:

- **Expedited Forwarding:** packet departure rate of a class equals or exceeds specified rate
  - logical link with a minimum guaranteed rate
- **Assured Forwarding:** e.g. 4 classes of traffic
  - each class guaranteed minimum amount of bandwidth and a minimum of buffering
  - packets each class have one of three possible drop preferences; in case of congestion routers discard packets based on drop preference values



# Assured Forwarding DiffServ Code Points

- Assured Forwarding behavior definition
  - RFC 2597 - Juha Heinanen, Fred Baker, Walter Weiss, John Wroclawski: Assured Forwarding PHB Group
    - Recommended Codepoints: c.f. table below
  - RFC 3260 - Dan Grossman:  
New Terminology and Clarifications for Diffserv

## RFC 2597 Assured Forwarding (AF) Recommended Code Points

	<b>Class 1 (lowest)</b>	<b>Class 2</b>	<b>Class 3</b>	<b>Class 4 (highest)</b>
<b>Low Drop</b>	001010	010010	011010	100010
<b>Med Drop</b>	001100	010100	011100	100100
<b>High Drop</b>	001110	010110	011110	100110