

Network Security

Chapter 1

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Introduction

Network InSecurity

- ▶ By example: An Ethernet cable
- ▶ How secure is it?



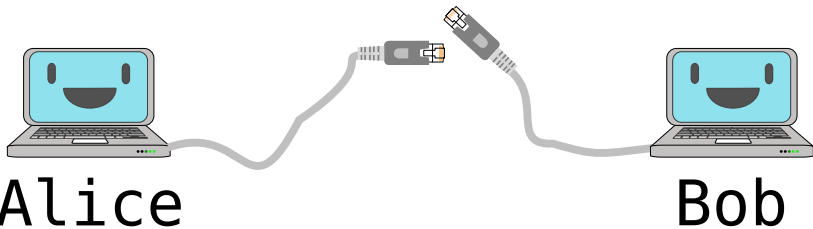
Alice



Bob

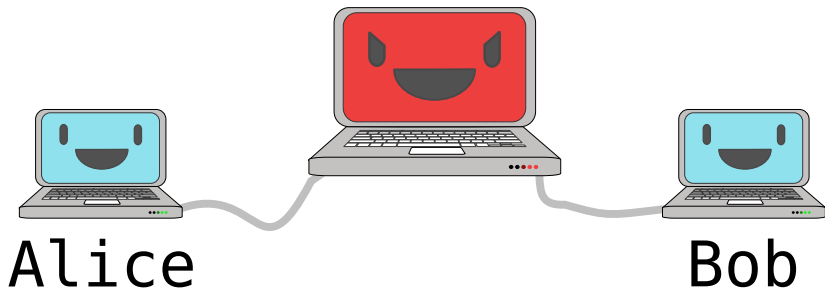
Network InSecurity

- ▶ Step 1: Obtain a knife
- ▶ Step 2: Add RJ45 adapters



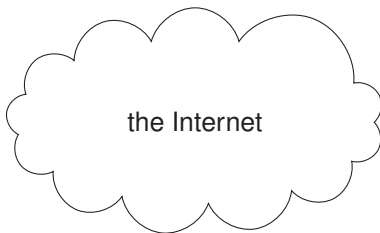
Network InSecurity

- ▶ Step 3: Configure transparent ethernet bridging
- ▶ You are now in full control of the traffic
 - ▶ read
 - ▶ modify
- ▶ Technical term: *Man in the Middle* (MitM)



Network InSecurity

Alice



Bob

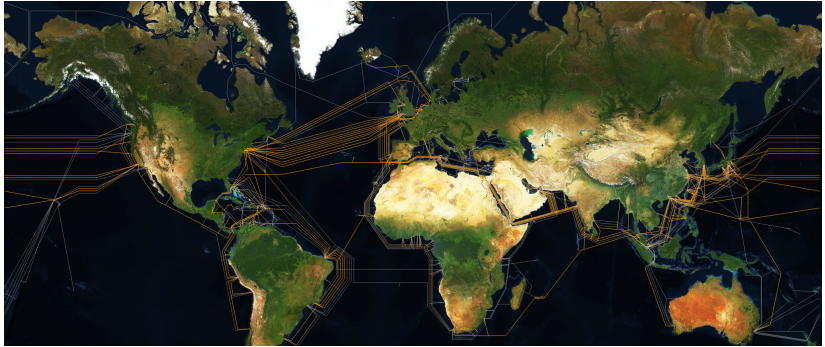
Network InSecurity

Alice



Bob

Network “Security” offered by our Secret Services



<http://lifewinning.com/submarine-cable-taps/>

- ▶ Passive attacks: wiretapping, ...
- ▶ Active attacks: Quantum Insert, ...
- ▶ Combined: economic espionage, ...



Attackers

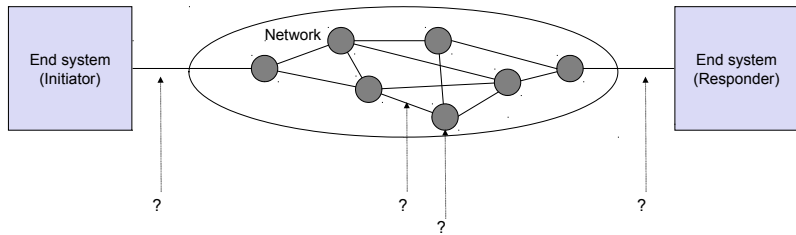
Attacker Models

- ▶ Attacking communications on the message level
- ▶ Passive attacks:
 - ▶ Eavesdropping of messages
- ▶ Active attacks
 - ▶ all passive attacks
 - ▶ Delay
 - ▶ Replay
 - ▶ Deletion
 - ▶ Modification
 - ▶ Insertion

General Attacker Model

- ▶ The attacker *is* the network
- ▶ And can perform any active attack
- ▶ But cannot break cryptographic primitives
- ▶ This is called the Dolev-Yao attacker model
- ▶ If not stated otherwise, we will always assume this attacker model.

Attackers Limited by their Position in the Network



Attackers Limited by their Position in the Network

- ▶ Assume the Attacker is close to you
- ▶ Example: You sit in a cyber cafe and accidentally connected to the attacker's hotspot
 - ▶ The attacker can perform any active attacks on you
 - ▶ But you can bypass this attacker: Establish a *secure* tunnel to a server in the Internet
 - ▶ Route all your packets over the secure tunnel
 - ▶ The attacker can now perform only DOS (Denial Of Service) attacks against you

Attackers Limited by their Position in the Network

- ▶ Assume the Attacker is close to your servers
- ▶ Example: She rented a VM on the same host machine where your virtual server is running
 - ▶ The attacker could try to perform timing attacks against you
 - ▶ By measuring how long certain operations take at your server, the attacker might be able to break a security service
 - ▶ (only if the service is vulnerable to side channel attacks)
 - ▶ Such measurement is usually not possible over the Internet

Attackers Limited by their Position in the Network

- ▶ Assume the Attacker is somewhere in the Internet
- ▶ Internet: Best effort packet switching
- ▶ End-user has no control how packets are routed
- ▶ Are all AS/ISP trustworthy?
- ▶ Does your ISP alter your packets?
 - ▶ “value added service” i.e. your ISP places advertisement on the websites you are visiting
- ▶ BND/NSA/GCHQ black boxes are basically everywhere

Security Goals

Security Goals Technically Defined

- ▶ Data Integrity
 - ▶ No improper or unauthorized change of data
- ▶ Confidentiality
 - ▶ Concealment of information
- ▶ Availability
 - ▶ Services should be available and function correctly
- ▶ Authenticity
 - ▶ Entity is who she claims to be
- ▶ Accountability german: „Zurechenbarkeit“
 - ▶ Identify the entity responsible for any communication event
- ▶ Controlled Access
 - ▶ Only authorized entities can access certain services or information

Quiz

- ▶ What is needed to support non-repudiation?
(„*Nicht-Abstreitbarkeit*“)

Quiz

- ▶ What is needed to support non-repudiation?
(„*Nicht-Abstreitbarkeit*“)
 - ▶ Accountability

Quiz

- ▶ What is necessary to support accountability?

Quiz

- ▶ What is necessary to support accountability?
 - ▶ Authenticity

Quiz

- ▶ What do you want to support deterrence („*Abschreckung*“)

Quiz

- ▶ What do you want to support deterrence („*Abschreckung*“)
 - ▶ Accountability

Quiz

- ▶ What is data origin integrity?

Quiz

- ▶ What is data origin integrity?
 - ▶ Authenticity

Quiz

- ▶ What is the difference?
- ▶ Authentication

- ▶ Authorization

Quiz

- ▶ What is the difference?
- ▶ Authentication
 - ▶ Proves who you are
 - ▶ Associated security goal: Authenticity
- ▶ Authorization
 - ▶ Defines what you are allowed to do
 - ▶ Associated security goal: Controlled Access

Quiz

- ▶ What is the difference?
- ▶ Authentication
 - ▶ Proves who you are
 - ▶ Associated security goal: Authenticity
 - ▶ E.g. your passport
- ▶ Authorization
 - ▶ Defines what you are allowed to do
 - ▶ Associated security goal: Controlled Access
 - ▶ E.g. “*are you on the VIP list?*”

Mixing Authentication and Authorization



My best attempt was registering to Black Hat with first name: "Staff" and last name: "Access All Areas"

<https://twitter.com/mikko/status/587973545797492738>



Threats

Threats

- ▶ Abstract Definition
 - ▶ A threat in a communication network is any possible event or sequence of actions that might lead to a violation of one or more security goals
 - ▶ The actual realization of a threat is called an attack

Threats Technically Defined

- ▶ Masquerade
 - ▶ An entity claims to be another entity (also called “impersonation”)
- ▶ Eavesdropping
 - ▶ An entity reads information it is not intended to read
- ▶ Loss or Modification of (transmitted) Information
 - ▶ Data is being altered or destroyed
- ▶ Denial of Communication Acts (Repudiation)
 - ▶ An entity falsely denies its participation in a communication act
- ▶ Forgery of Information
 - ▶ An entity creates new information in the name of another entity
- ▶ Sabotage/Denial of Service
 - ▶ Any action that aims to reduce the availability and / or correct functioning of services or systems
- ▶ Authorization Violation:
 - ▶ An entity uses a service or resources it is not intended to use

Example 1

- ▶ Eavesdropping + Authorization Violation
- ▶ Example
 - ▶ Alice@Box\$./rootremoteshell \$ROUTER
root@router# tcpdump | grep password
- ▶ If Alice does not start modifying the traffic, she is a passive attacker
- ▶ Note: If not stated otherwise, we assume that attackers don't have remote code execution on our boxes

Example 2

- ▶ Masquerade + Forgery of Information
- ▶ Example
 - ▶ Alice pretends to be Bob
 - ▶ Alice@Box\$ `hping3 --count 1 --spooof $BOB --icmp --icmptype 8 $CARL`
 - ▶ Bob gets an ICMP Echo Reply which he never requested
- ▶ Alice is an active attacker

Example 2: IP Spoofing cont.

Alice

Bob

Carl

A line with an arrow pointing from Alice to Carl, containing the text: `src:Bob dst:Carl ping`

A line with an arrow pointing from Carl to Bob, containing the text: `src:Carl dst:Bob pong`

Example 2: IP Spoofing cont.

- ▶ Alice: 192.168.1.170
- ▶ Bob 192.168.1.227
- ▶ Carl: 192.168.1.1
- ▶ Alice sends the spoofed packet
 - ▶ Internet Protocol Version 4, Src: **192.168.1.227**, Dst: 192.168.1.1; ICMP Echo Request
- ▶ Carl replies to the source address specified
- ▶ Bob receives a lonely echo reply
 - ▶ Internet Protocol Version 4, Src: 192.168.1.1, Dst: 192.168.1.227; ICMP Echo Reply

192.168.1.1 192.168.1.227 ICMP 60 Echo (ping) reply id=0xce1f, seq=0/0, ttl=61

Example 3

- ▶ Denial of Service
- ▶ Example
 - ▶ Bob runs a webserver (http, tcp port 80) with very few memory
 - ▶ Alice floods Bob with TCP SYN packets
 - ▶ Alice@Box\$ `hping3 --fast --count 42 --syn --destport 80 $BOB`
 - ▶ Bob allocates memory to store the 42 connections in the SYN-RECEIVED state
- ▶ Now Alice starts to deny that she is responsible for the attack
- ▶ Denial of Service + Forgery of Information + Denial of Communication Acts
- ▶ Example
 - ▶ Alice@Box\$ `hping3 --fast --count 42 --rand-source --syn --destport 80 $BOB`
 - ▶ `--rand-source`: random spoofed source IP address

Example 3

Capturing from Ethernet [Wireshark 1.12.4 (v1.12.4-0-gb4861da from master-1.12)]

File Edit View Go Capture Analyze Statistics Telephony Tools Internals Help

Filter: Expression... Clear Apply Save

No.	Time	Source	Destination	Protocol	Length	Info
736	686.042764000	56.10.51.117	192.168.1.227	TCP	60	1350-80 [SYN] Seq=0 win=512 Len=0
737	686.129344000	38.36.23.85	192.168.1.227	TCP	60	1351-80 [SYN] Seq=0 win=512 Len=0
738	686.229507000	36.116.117.78	192.168.1.227	TCP	60	1352-80 [SYN] Seq=0 win=512 Len=0
739	686.329714000	189.139.51.172	192.168.1.227	TCP	60	1353-80 [SYN] Seq=0 win=512 Len=0
740	686.429848000	242.114.151.137	192.168.1.227	TCP	60	1354-80 [SYN] Seq=0 win=512 Len=0
741	686.530802000	255.124.118.119	192.168.1.227	TCP	60	1355-80 [SYN] Seq=0 win=512 Len=0
742	686.630208000	161.10.181.62	192.168.1.227	TCP	60	1356-80 [SYN] Seq=0 win=512 Len=0
743	686.730401000	9.205.193.205	192.168.1.227	TCP	60	1357-80 [SYN] Seq=0 win=512 Len=0
744	686.830479000	205.95.119.125	192.168.1.227	TCP	60	1358-80 [SYN] Seq=0 win=512 Len=0
745	686.930632000	238.97.119.210	192.168.1.227	TCP	60	1359-80 [SYN] Seq=0 win=512 Len=0
746	687.030809000	194.238.30.56	192.168.1.227	TCP	60	1360-80 [SYN] Seq=0 win=512 Len=0
747	687.130950000	111.148.162.200	192.168.1.227	TCP	60	1361-80 [SYN] Seq=0 win=512 Len=0
748	687.230995000	225.60.95.186	192.168.1.227	TCP	60	1362-80 [SYN] Seq=0 win=512 Len=0
749	687.331114000	124.161.110.246	192.168.1.227	TCP	60	1363-80 [SYN] Seq=0 win=512 Len=0
750	687.431808000	193.202.206.237	192.168.1.227	TCP	60	1364-80 [SYN] Seq=0 win=512 Len=0

- ▶ Why does the attack succeed?
- ▶ This is a good opportunity to refresh your knowledge about the TCP 3-way handshake

Literature

- ▶ Matt Bishop, *Introduction to Computer Security*, Addison-Wesley, 2004
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