## Network Security

## Chapter 2 Basics of Cryptography

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- Overview of Cryptographic Algorithms <br> - Attacking Cryptographic Algorithms <br> - Historical Approaches
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- Foundations of Modern Cryptography


## Cryptographic algorithms: overview

- During this course two main applications of cryptographic algorithms are of principal interest:
- Encryption of data: transforms plaintext data into ciphertext in order to conceal its' meaning
- Signing of data: computes a check value or digital signature to a given plain- or ciphertext, that can be verified by some or all entities being able to access the signed data
- Some cryptographic algorithms can be used for both purposes, some are only secure and / or efficient for one of them.
- Principal categories of cryptographic algorithms:
- Symmetric cryptography using 1 key for en-/decryption or signing/checking
- Asymmetric cryptography using 2 different keys for en-/decryption or signing/checking
- Cryptographic hash functions using 0 keys (the "key" is not a separate input but "appended" to or "mixed" with the data).


## Attacking cryptography (1): Cryptanalysis

- Cryptanalysis is the process of attempting to discover the plaintext and / or the key
- Types of cryptanalysis:
- Ciphertext only: specific patterns of the plaintext may remain in the ciphertext (frequencies of letters, digraphs, etc.)
- Known ciphertext / plaintext pairs
- Chosen plaintext or chosen ciphertext
- Newer developments: differential cryptanalysis, linear cryptanalysis
- Cryptanalysis of public key cryptography:
- The fact that one key is publicly exposed may be exploited
- Public key cryptanalysis is more aimed at breaking the cryptosystem itself and is closer to pure mathematical research than to classical cryptanalysis
- Important directions:
- Computation of discrete logarithms
- Factorization of large integers


## Attacking cryptography (2): brute force attack

- The brute force attack tries every possible key until it finds an intelligible plaintext:
- Every cryptographic algorithm can in theory be attacked by brute force
- On average, half of all possible keys will have to be tried

|  | Average Time Required for Exhaustive Key Search |  |  |
| :---: | :--- | :--- | :--- |
| Key Size [bit] | Number of keys | Time required <br> at 1 encryption $/ \mu \mathrm{s}$ | Time required <br> at $10^{6}$ encryption/ $\mu \mathrm{s}$ |
| 32 | $2^{32}=4.3 * 10^{9}$ | $2^{31} \mu \mathrm{~s}=35.8$ minutes | 2.15 milliseconds |
| 56 | $2^{56}=7.2 * 10^{16}$ | $2^{55} \mu \mathrm{~s}=1142$ years | 10.01 hours |
| 128 | $2^{128}=3.4 * 10^{38}$ | $2^{127} \mu \mathrm{~s}=5.4 * 10^{24}$ years | $5.4 * 10^{18}$ years |

- 1 encryption / $\mu \mathrm{s}$ : 100 Clock cycles of a 100 MHz processor
- $10^{\wedge} 6$ encryptions / $\mu \mathrm{s}$ : Clock cycles using 500 parallel 2 GHz processors

Attacking cryptography (3): How large is large?

Reference Numbers Comparing Relative Magnitudes

| Reference Numbers Comparing Relative Magnitudes |  |  |
| :--- | :--- | :--- |
| Reference | Magnitude |  |
| Seconds in a year | $\approx 3 \quad * 10^{7}$ |  |
| Seconds since creation of solar system | $\approx 2$ | $* 10^{17}$ |
| Clock cycles per year (3 GHz computer) | $\approx 1 \quad * 10^{17}$ |  |
| Binary strings of length 64 | $2^{64}$ | $\approx 1.8 * 10^{19}$ |
| Binary strings of length 128 | $2^{128}$ | $\approx 3.4 * 10^{38}$ |
| Binary strings of length 256 | $2^{256}$ | $\approx 1.2 * 10^{77}$ |
| Number of 75 -digit prime numbers | $\approx 5.2 * 10^{72}$ |  |
| Electrons in the universe | $\approx 8.37 * 10^{77}$ |  |

## Classification of modern encryption algorithms

- The type of operations used for transforming plaintext to ciphertext:
- Substitution, which maps each element in the plaintext (bit, letter, group of bits or letters) into another element
- Transposition, which re-arranges elements in the plaintext
- The number of keys used:
- Symmetric ciphers, which use the same key for en- / decryption
- Asymmetric ciphers, which use different keys for en- / decryption
- The way in which the plaintext is processed:
- Stream ciphers work on bit streams and encrypt one bit after another
- Block ciphers work on blocks of width $b$ with $b$ depending on the specific algorithm.


## Basic Kryptographic Principles

- Substitution
- Individual characters are exchanged by other characters

Types of substitution

- simple substitution substitution: operates on single letters
- polygraphic substitution: operates on larger groups of letters
- monoalphabetic substitution: uses fixed substitution over the entire message
- polyalphabetic substitution: uses different substitutions at different sections of a message
- Transposition
- The position of individual characters changes (Permutation)


## Transposition: scytale

- Known as early as $7^{\text {th }}$ century BC
- Principle:
- Wrap parchment strip over a wooden rod of a fixed diameter and write letters along the rod.
- Unwrap a strip and "transmit"
- To decrypt, wrap a received over a wooden rod of the same diameter and read off the text.
- Example:
troops
headii
nthewe
stneed $\Rightarrow \quad$ thnsm predd opoah nrlod eeeis iedus
stneed
pplies
- Weakness:
- Easy to break by finding a suitable matrix transposition.


## Monoalphabetic substitution: Atbash

Jeremiah 25:25
And all the kings of the north, far and near, one with another, and all the kingdoms of the world, which are upon the face of the earth: and the king of Sheshach shall drink after them.

Atbash code: reversed Hebrew alphabet.

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| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $\begin{array}{\|c\|} \hline \frac{\top}{\top} \\ \frac{\mathrm{Taw}}{n} \end{array}$ | $\begin{aligned} & \hline s \\ & \frac{s i n}{v} \end{aligned}$ | $\begin{array}{\|c\|} \hline R \\ \hline \frac{R}{R} \\ \hline \end{array}$ | $\begin{array}{\|c} \hline \frac{Q}{\mathrm{Koph}} \\ \hline \end{array}$ | $\begin{array}{\|c\|} \hline \frac{\mathrm{Sade}}{} \\ \hline \mathrm{~s} \end{array}$ | $\begin{aligned} & \hline p \\ & \frac{P e}{72} \end{aligned}$ | $\begin{array}{\|c} \hline \frac{O}{\text { Ainn }} \end{array}$ | $\begin{array}{\|c\|} \hline x \\ \frac{\text { Samech }}{0} \\ \hline \end{array}$ | $\begin{array}{\|c\|} \hline \frac{N}{\left\lvert\, \frac{N u n}{12}\right.} \end{array}$ | $\begin{array}{\|c\|c\|} \hline \left.\begin{array}{c} \mathrm{m} \\ \mathrm{Mem} \\ \hline 0 \end{array} \right\rvert\, \end{array}$ | $\begin{array}{\|c\|} \hline \frac{\mathrm{Lamed}}{i} \\ \hline \end{array}$ | $\begin{gathered} \hline K \\ \frac{K a p h}{7} \end{gathered}$ | $\begin{aligned} & \text { Jod } \\ & \text { Jood } \end{aligned}$ | $\begin{array}{\|c\|} \hline \frac{T}{4} \\ \frac{T e t}{v} \end{array}$ | $\begin{aligned} & \begin{array}{c} \text { H } \\ \text { Chet } \end{array} \end{aligned}$ | $\begin{array}{\|c\|} \hline \text { Zajin } \end{array}$ | wVFY Waw | $\begin{aligned} & \hline \mathrm{H} \\ & \underline{\mathrm{He}} \mathrm{~J} \end{aligned}$ | D ${ }_{\text {D }}^{\text {Datet }}$ | ${ }_{\text {Gimel }}^{\text {G }}$ | $\frac{\text { Beth }}{}$ | $\underset{\text { A }}{\text { Aleph }}$ |



Monoalphabetic substitution: Caesar cipher

- Ceasar code: left shift of alphabet by 3 positions.

- Example (letter of Cicero to Caesar): MDEHV RSNQNRQNV PHDH XHVXNPRQNZP HABES OPINIONIS MEAE TESTIMONIUM
- Weakness: a limited number of possible substitutions. Easy to break by brute force!


## Modern cryptography: S and P-boxes

S-box:

- Block-wise substitution of binary digits.
- Resistant to attacks for sufficiently large block size; e.g. for $\mathrm{n}=128$ it provides $2^{128}$ possible mappings.


P-box:

- Block-wise permutation of binary digits.
- Realizes a simple transposition cipher with maximal entropy.
- Problem: straightforward attacks exist.



## A practical Feistel cipher

- A multiple-round scheme with separate keys per round.
- Invertible via a reverse order of rounds.
- 3 rounds suffice to achieve a pseudorandom permutation.
- 4 rounds suffice to achieve a strong pseudorandom permutation (i.e. it remains pseudorandom to an attacker with an oracle access to its inverse permutation).
- A foundation for a large number of modern symmetric ciphers: DES, Lucifer, Blowfish, RC5, Twofish, etc.



## Important properties of encryption algorithms

Consider, a sender is encrypting plaintext messages $\mathrm{P}_{1}, \mathrm{P}_{2}, \ldots$ to ciphertext messages $\mathrm{C}_{1}, \mathrm{C}_{2}, \ldots$
Then the following properties of the encryption algorithm are of special interest:

- Error propagation characterizes the effects of bit-errors during transmission of ciphertext to reconstructed plaintext $\mathrm{P}_{1}{ }^{\prime}, \mathrm{P}_{2}{ }^{\prime}, \ldots$
- Depending on the encryption algorithm there may be one or more erroneous bits in the reconstructed plaintext per erroneous ciphertext bit
- Synchronization characterizes the effects of lost ciphertext data units to the reconstructed plaintext
- Some encryption algorithms can not recover from lost ciphertext and need therefore explicit re-synchronization in case of lost messages
- Other algorithms do automatically re-synchronize after 0 to n ( n depending on the algorithm) ciphertext bits

